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UTILITY PATENT APPLICATION TRANSMITTAL <small>(Only for new nonprovisional applications under 37 C.F.R. § 1.53(b))</small>	Attorney Docket No. _____
	First Inventor or Application Identifier <u>Gerd Krämer</u>
	Title <u>(Automatic Programming)</u>
	Express Mail Label No. _____

APPLICATION ELEMENTS <small>See MPEP chapter 600 concerning utility patent application contents</small>	ADDRESS TO: Assistant Commissioner for Patents Box Patent Application Washington, DC 20231
1. <input checked="" type="checkbox"/> * Fee Transmittal Form (e.g., PTO/SB/17) <small>(Submit an original and a duplicate for fee processing)</small> 2. <input checked="" type="checkbox"/> Specification [Total Pages <u>32</u>] <small>(preferred arrangement set forth below)</small> - Descriptive title of the Invention - Cross References to Related Applications - Statement Regarding Fed sponsored R & D - Reference to Microfiche Appendix - Background of the Invention - Brief Summary of the Invention - Brief Description of the Drawings (if filed) - Detailed Description - Claim(s) - Abstract of the Disclosure 3. <input checked="" type="checkbox"/> Drawing(s) (35 U.S.C. 113) [Total Sheets <u>19</u>] 4. Oath or Declaration [Total Pages <u>3</u>] a. <input checked="" type="checkbox"/> Newly executed (original or copy) b. <input type="checkbox"/> Copy from a prior application (37 C.F.R. § 1.53(b)(2)) i. <input type="checkbox"/> DELETION OF INVENTOR(S) Signed statement attached deleting inventor(s) named in the prior application, see 37 C.F.R. §§ 1.53(b)(2) and 1.53(b)(3)	5. <input type="checkbox"/> Microfiche Computer Program (Appendix) 6. Nucleotide and/or Amino Acid Sequence Submission (if applicable, all necessary) a. <input type="checkbox"/> Computer Readable Copy b. <input type="checkbox"/> Paper Copy (identical to computer copy) c. <input type="checkbox"/> Statement verifying identity of above copies ACCOMPANYING APPLICATION PARTS 7. <input type="checkbox"/> Assignment Papers (cover sheet & document(s)) 8. <input type="checkbox"/> 37 C.F.R. § 1.53(b) Statement of Power of Attorney (when there is an assignee) 9. <input checked="" type="checkbox"/> English Translation Document (if applicable) 10. <input type="checkbox"/> Information Disclosure Statement (IDS)/PTO-1449 <input type="checkbox"/> Copies of IDS Citations 11. <input type="checkbox"/> Preliminary Amendment 12. <input type="checkbox"/> Return Receipt Postcard (MPEP 503) (Should be specifically Remized) 13. <input type="checkbox"/> Small Entity Statement <input type="checkbox"/> Statement filed in prior application, Status still proper and desired 14. <input checked="" type="checkbox"/> Certified Copy of Priority Document(s) <u>2</u> (if foreign priority is claimed) 15. <input type="checkbox"/> Other: <u>Application Data Sheet</u>
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Examiner Name
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Attorney Docket No. **1**

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FEE CALCULATION

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101 710 201 355 Utility filing fee	345
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107 480 207 245 Plant filing fee	1
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Large Entity Small Entity Fee Fee Fee Code (\$) Code (\$) Fee Description	
103 18 203 9 Claims in excess of 20	
102 80 202 40 Independent claims in excess of 3	
104 270 204 135 Multiple dependent claim, if not paid	
109 80 209 40 ** Reissue independent claims over original patent	
110 18 210 9 ** Reissue claims in excess of 20 and over original patent	

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139 130 139 130 Non-English specification	
127 2,520 147 2,520 For filing a request for ex parte reexamination	
112 920* 112 920* Requesting publication of SIR prior to Examiner action	
113 1,840* 113 1,840* Requesting publication of SIR after Examiner action	
115 110 215 55 Extension for reply within first month	
116 390 216 195 Extension for reply within second month	
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142 1,240 242 620 Utility issue fee (or reissue)	
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149 710 249 355 For each additional invention to be examined (37 CFR § 1.129(b))	
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169 900 169 900 Request for expedited examination of a design application	

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APPLICATION INFORMATION

Title Line One:: Method for generating a simple kind of
Title Line Two:: Artificial Consciousness in a computer,
Title Line Three:: which has the capability to plan,
Title Line Four:: generate automatically and execute
Title Line Five:: machine-code for the solution of
Title Line Six:: arbitrary programming-abandonments.
Title Line Seven:: (Automatic Programming)
Total Drawing Sheets:: 19
Formal Drawings?:: Yes
Application Type:: Utility

PRIOR FOREIGN APPLICATIONS

Foreign Application One:: 199 52 587.0-53
Filing Date:: 02. November 1999
Country:: Germany
Priority Claimed:: Yes

09704807-40300

TITLE OF THE INVENTION:

Method for generating a simple kind of Artificial Consciousness in a computer, which has the capability to plan, generate automatically and execute machine-code for the solution of arbitrary programming-abandonments. (Automatic Programming)

REFERECES:

First announcement of this utility-patent "Verfahren zur Generierung einer einfachen Form künstlichen Bewußtseins im Computer zur Befähigung selbsttätig planender Erstellung von Maschinencode-Programmen und deren Ausführung zur Lösung beliebiger gestellter Programmieraufgaben" was the 2nd November 1999 in Germany at the DPMA (deutsches Patent- und Markenamt = german Patent- and Trademark-Office).

SPONSORSHIP STATEMENT:

There was no sponsor for this invention and I'm a single independent inventor.

BACKGROUND OF THE INVENTION:

1. Field of the Invention:

The present invention relates to computer programming, and more particularly to automatic code generation. It relates also to learn-capable programs and artificial intelligence.

2. Present State of the Art:

Worldwide in the Field of Software-Development many employees are missing and the development tasks become larger and larger.

Until now a given conceptual formulation is conceived and programmed by Software-developers.

For relieving the programming there are "Wizards" which offer the possibility to generate basic parts of source-code after making interactive inputs on dialog-windows by a fixed given generator-scheme.

Moreover company-specific scripts are written, which generate simple steadily repeated parts of source-code with variations on the same positions by reading out data out of ASCII-files.

In every case the user first has to develop the generating script and then has to write the ASCII data for to read out, or - in the case of "Wizards" - has to make user defined inputs and after the generation of the frame-sourcecode has to develop the intrinsic functionality of the program. After it the source-code has to be compiled to become executable. But such programs are not adaptive.

On the area of AI there are neuronal networks / fuzzy logic which can build expert-systems, which can absorb external attractions and have the capability to make adaptive decisions on these inputs, which means a kind of adaptive control system but they will not be able to plan and develop and execute machine-code and learn from its execution.

In the decision 20 W (pat) 12/76 of the german patent court artificial consciousness was tried to generated in a patent application by a reflexive chain of video cameras and monitors - this procedure has nothing to do with that method.

BRIEF SUMMARY OF THE INVENTION:

It's an object of the invention to create computer based artificial consciousness.

It is another object of the present invention to provide a method for giving a computer the capability to learn programming for itself.

It's a further object of the invention to provide a system to make a computer plan and develop programs targeted to a pregiven programming-aim or to fulfill its basic needs.

Therefore it first captures all processor exception-vectors (for a single process-system) or task exception-vectors (for a multitasking version) by own analysis-routines.

Then the system generates numbers, puts them to a defined place in memory and then sets the instruction-pointer to that number for to execute it, like it would be a legal opcode.

Before the number is executed the processor's registers are set to predefined initial conditions and one number is executed several times using different initial conditions.

After every number-execution the system analyses, if an exception occurred or the number caused a jump or a write to memory and the concerned source- and destination-registers are determined and also the kind of instruction, which means its mnemonic. For every execution many theoretical source-registers and several possible commands are possible. Therefore one number is executed by

many predefined different initial conditions to determine the concerned source-register and operation most exactly. By the sum of the execution analysis data the command concerned mnemonic and the source- and destination-registers are determined. So the system itself learns to program in machine-code.

Additional the absolute basic needs, which also have mono-cellars, are modelled:

Pain means an attack to the program (=overwrite) and hunger means loss of energy, which is modelled by a defined register (hunger=low values).

The program has got two valuation-systems: one concerning the basic needs and one concerning the fulfillment of pregiven programming aim.

After single numbers are executed two-number-combinations are executed and the effects of these combinations are determined.

The valuation system determines if the combination is good or worse concerning the basic needs and the fulfillment of a pregiven programming-aim (it's possible to disable one of these two valuation-systems).

The programming-aim concerning valuation-function is dynamic, which means the value-range of its valuation-results is valued by a meta-valuation-system - and if the valuation-results are not very meaningful, which means, they're clustered near the min/max-boundaries or near zero or another value, the valuation-function is changed by the valuation-system itself and a revaluation occurs. If then the valuations results are worse than before, the modification is quashed and another valuation-function modification is tried until the revaluation results in unclustered valuation-results.

When larger number-combinations are tried, the valuation-system omits to combine combinations which caused fatal exceptions, large jumps, extensive writing to memory, registers which should not be used, etc. or combinations which dislodge from the programming-aim. So not every additional number or combination in the total combination causes an exponential rise of needed calculation-time and disc-space.

The larger the combination becomes the more restrictive is the programming-aim specific valuation-system concerning additional numbers=opcodes or combinations. Then additional combinations must appropriate the programming-aim.

So the system learns to plan developing the desired routine.

The solution-routines are retested by nearly all possible input-values and if it works fine it's

valuated by needed clock cycles and memory space and the most effective solution-routine then is disassembled and can be taken by developers to implement it into their projects as a subroutine.

BRIEF DESCRIPTION OF THE DRAWINGS:

Fig.1 shows the ER-diagram of the database-tables which contain the data of the AC-program. The in the middle shown CLT(i)-tables are created dynamically. The database is described in section 1.3.2.

Fig.2-18 show the names, datatypes, value-range and meanings of the table's columns - some with additional examples, how they're filled. These tables are described in sections 1.3.2.1 to 1.3.2.16.

Fig.19-21 show the value-assignments of the O_xT and C_xT-tables. Here the effects of the executions are analyzed, and the mnemonic and source+destination-registers are determined.

Fig.22 shows the value-assignments of the energyspecific tables ELT and EBT, which provide the analysis results of the actions concerning the modelled hunger.

Fig.23-24 show the flow-chart of the AC-program.

DETAILED DESCRIPTION OF THE INVENTION:

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1. SPECIFICATION

1.1 Object of the invention:

The objects of the invention are:

- a) to provide automatized software-development,
- b) to create computerbased artificial consciousness.

With this method a simple form of artificial consciousness is generated using a computer, which acts aimless and arbitrarily in the beginning, but has the capability to learn from the effects of all its "behaviour", for to, when it knows the effects of every single behaviour or behaviour pattern (=combination), has the capability to join together its single actions targeting to a given aim or fulfilment of basic needs.

1.2 Derivation of technical feasibility and Definition of artificial Consciousness:

1.2.1 Philosophical basic liberations:

(... are normally no ingredient of a patent specification, but indispensable for the explanation of the technical feasibility)

If the prerequisite of consciousness of man would be a kind of soul which would be adventitious between cygot and birth, it would be possible to locate it by cogitation experiment:

If you would cognitive cut the head and provide the carotids with oxygen and nutrient containing blood, the consciousness would surely be located in the head.

If you would isolate the brain expect of the factitious supply, no conventionally information flow between the individual and the environment would be possible, but the "I am"-consciousness would surely be present.

Over it it's theoretical now possible to cut the cerebral lappets for seeing, listening, smelling, tasting, feeling, equilibration, speech and the cerebellum and nothing more would be lost.

If the front cerebral lappets would further be cut, you'd loose the possibility to compute with present knowledge and on the cut of several upper cerebral lappets you'd loose reminiscence, but the deepest basic "I am" would be remaining.

⇒ If a kind of soul would exist, it would be located on the upper End of the phylum brain ##.

Under consideration of the fluent evolution the primates would have a "soul" too; and other mammals too; and all other animals too; and the single-cellars too; and plants too; and consequently every cell of a multicellular life form too.

⇒ A border of a "soul-adventitious" or an "I am"-consciousness between the life forms is

missing.

- ⇒ Every cell of our body ought have its own soul (the evolutionary specialisation to neural-cells is, equivalent to the prenatal cell-fissions, fluent).
- ⇒ A soul does not exist (it's not necessary to build consciousness).
- ⇒ Consciousness originates during the evolution compulsively automatically.
- ⇒ Inside the "dead" molecule of the DNA is the construction plan for generating consciousness.
- ⇒ Consciousness originates by valuation of the own actions and its effects, with the reflection of the valuation-results on the adapting dynamic valuation-method.
- ⇒ If no soul is necessary for generating consciousness, but only the complex "program" of DNA, then consciousness is also generatable by a complex reflexive computer-program.

1.2.2 Basic Approach for the Generation of Artificial Consciousness:

The doing of all, including the plainest individuals conduces the fruition of its basic needs.

These basic needs are:

- a) no achiness := no attack against the own system *and*
- b) no hunger := no imminent loss of energy

A complex program, in which these basic needs are modelled, and which can act freely and has the possibility to learn reflexively what its actions effectuate (like a child) and can reflect if its actions improves its situation in reference to its basic needs, builds up a valuation-system, and then plans with the actions from its learned knowledge and reflects again and so attains consciousness. When it later scans its own machine-code and then tests every of its opcodes and after it all opcode-combinations it discerns the effect of its opcode-combinations and their context and so attains self-consciousness and now has the possibility of self-reproduction and the aware improvement of its machine-code while reproduction and so starts its evolution (its so effective like man could improve its DNA while mitoses/meiosis implementing its experience of life).

1.3. Technical Doctrine how to generate Artificial Consciousness:

1.3.0 Definition of the appropriated abbreviations:

The programming works on every computer with any processor and on any operating-system. In the following μ -indexed abbreviations correlate with Motorola processors, π means Intel processors, and ψ -indexed denote PowerPC-RISC-Processors.

eq. = equivalent

ASR_ψ = Address Space Register

BAT_ψ = BAT-Registers

BC = BitCode: every Bit correlates with a Flag and combinations are allowed.

CCR_{xi} = Condition-Code-Register (= Flags: EXtension, Negative-, Zero-, Overflow-, Carry-)

CISC = Complex-InstructionSet Computer (p.e. IA₃₂ and MC₆₈)

CPU = Central processing unit = Prozessor

CR_ψ = Condition-Register (CR 0..7)

CR_π = Control-Register

CTR_ψ = Count-Register

DABR_ψ = Data Address Breakpoint Register

DAR_ψ = Data Address Register

DB = DataBase

DEC_ψ = Decrement-Register

DR_π = Debug-Register

DSISR_ψ = DSI Status-Register shows the reason for DSI- and Alignment-Exceptions.

EA = Effective Address (memory access without using a register)

EAR_ψ = External Access Register

Reg_{στ} = Segment Register: CS; SS; DS, ES, FS, GS

EFlags_π = 32-Bit-Register with the System-Flags: Ident-, VirtualInterruptPending-, VirtualInterruptFlag-, AlignmentCheck-, Virtual8086Mode-, ResumeFlag-, NestedTask, InputOutputPrivilegeLevel, OverflowFlag, DirectionFlag, InterruptEnableFlag, TrapFlag, SignFlag, ZeroFlag, Auxiliary/Align-CarryFlag, ParityFlag, CarryFlag.

EIP_π = Extended Instruction Pointer (≙ PC_μ)

ER = Entity Relationship (database-model)

ESP_π = Extended StackPointer (≙ SP_μ)

Exc. = *Exception*_π: #DeviceError, #DeBug, NMI IRO, #BreakPoint, #OverFlow, #BoundRange exceeded, #UD (Invalid Opcode), #NM (device not available), #DoubleFault, invalid #TaskSwitch, Segment #NotPresent, #SS (StackFault), #GeneralProtection, #PageFault, #MF (FloatingPoint-Error), #AlignmentCheck, #MachineCheck.

*Exception*_μ: Reset, BusError, AddressError, invalidOpCode, Div/0, CHK, TrapV, PrivilegeViolation, Trace, Interrupts, Traps.

*Exception*_ψ: System-Reset, Machine-Check, DSI, ISI, Ext.Interrupt, Alignment, Program, Floating-Point unavailable, Decrementer, System Call, Trace, Floating-Point Assist.

FFT = fast Fourier-Transformation

FK = Foreign Key of a ER-Database-table

FPR_ψ = Floating-Point Register 0..31

FPSCR_ψ = Floating-Point Status and Control Register

GB = GigaBytes = 2³⁰ Bytes

IA = Intel-Architecture

AC = Artificial Consciousness

ld = Logarithm dualis = \log_2

MB = MegaBytes = 2^{20} Bytes

MSR_v = Machine State Register

NOP = NoOperation-OpCode [=opcode without effect (except increment of IP_{π}/PC_{μ})]

PC_{μ} = Program-Counter (=Pointer to the first byte in memory where the processor starts to fetch and execute an opcode)

PVR_{ψ} = Processor Version Register

RISC = Reduced-InstructionSet Computer (p.e. PowerPC₉₇₀)

RTE_μ = instruction: return from exception (loads SR and PC from Supervisor-Stack)

$$\text{SDR1}_{\psi} = \text{SDR1-Register}$$
$$\text{SPRG}_{\psi} = \text{SPRG } 0..3$$

SR_μ = StatusRegister (Flags_μ: Trace-, Supervisor-, + Interrupt-Maske: I₂ I₁ I₀, + **CCR**-Flags)

SR_π = Segment Registers: CS, DS, SS, ES, FS, GS

SR_{ψ} = Segment Registers

SRR_w = Save/Restore-Register of Machine-Status

SSP_μ = Supervisor-StackPointer (A7 in Supervisor-Modus)

TB _{π} = Time Base Facility: Time-Counter $\rightarrow 2^{64}-1$

TR_{*π*} = Table-Register: GDTR, IDTR, LDTR, TR

USP_μ = User-StackPointer (Adressregister A7 in User-Mode, A7' in Supervisor-Mode)

 $\hat{=}$ = correlates

\$ = begin of a hexadecimal number

ζ = result of bit-by-bit-AND over all following values [= $V_1 \& V_2 \& \dots \& V_n$]

$\text{OR} = \text{result of bit-by-bit-OR over all following values } [= V_1 \mid V_2 \mid \dots \mid V_n]$

r = number (sum) of the set Bits in the following value $[=(1 \& V) + (2 \& V)/2 + (4 \& V)/4 + \dots]$

\forall = for all of the following ...

 $\forall^* = \text{for all other ... (except the following)}$

1.3.1 Procedure for generating artificial consciousness by comprehensible words:

A computer system attains consciousness, if the active program, in which basic needs are modelled (see 1.3.5), has captured all exception-vectors and proceeds as follows:

Generate a normal number, write it somewhere in the memory, put the program-counter=instruction-pointer on it and execute it like it would be an opcode (=machine-code-command) and analyze, what its execution caused and proceed so with all numbers→opcodes (until a maximum length) with many representative initial conditions (register-values and reference-contents of address-registers).

Then use the saved opcodes which seldom caused an exception while using several initial conditions and evaluate if its execution increased or decreased its situation in due to its basic needs.

Combine the opcodes, which didn't decrease the own situation, and evaluate the effects of the code-combination using several initial conditions and save the result of analysis.

Plan combinations of those opcode-combinations which would increase the well being referable to the basic needs or could fulfill or approximate a given programming aim.

1.3.2 Creating the Database of the AC-Knowledge:

For to have the learned from actions persistent, and for to have convenient access to the large quantity of data, a relational database system with its tables and relations shown in 3.1 is created. For to increase access and to save hard disk capacity, equivalent primary keys in clusters and additional indexes for often used non-PK-rows. The ER diagram is shown in Fig.1.

Processor-dependant and in dependence of the number of 32-Bit-OpCodes, the database can grow very large and so the access speed according to it slow. Therefore RISC-Processors are more applicatively for the AC-Procedure than CISC-Processors. But CISC-Processors, like in the IA, use not very much opcodes which are longer than 24 bit, wherefore it's possible to have with striped tables and additional index-hard-disks and a higher "obliviousness" on inefficient opcode-combinations, an acceptable performance too.

The hard disc space problem in discussed in 1.5.

1.3.2.1 The Register-Identifikation-Table [RIT - Fig.2]:

In the RIT the individual descriptors of the processor-registers and -vectors are typed first: Every Register gets a correlated identification-number, an assigned bit in the BitCode, a character describing the register-type, a consecutive number of the register-type and an optional description of the register. The register of the processor-flags (EFlags_{pr}/SR_{pr}) gets the Register_ID #0. The exception-vectors mostly are located in memory and are no internal processor-registers - to identify most of these important vectors they get a Register_ID with negative sign, which correlates with

Fig.2b shows a Motorola example of the RIT.

Like in the RIT the registers, here in the OIT the most important operations get an identification number and a bit in a BitCode.

The *Operation_Type* pigeonholes the operation to an operation group, described in Fig.4c.

The *Operation_Mnemonic* (needn't to be exact like using assembler) and the optional *Operation_Description* describe the basic command.

Because of equal opcodes could cause different effects, in this table many representative initial conditions referring to all positive *Register_IDs* are pregiven here. For every initial condition number (in the fig.3-example: -31..+30) for all positive *Register_IDs* a sample of initial conditions is generated, p.e. using the Function in Fig.3b. But only for all registers, which could contain mathematically used numbers, like data-registers, address-register-references and in the decremented reference of it [to include the command `-(adr.reg.)`], floating-point-registers and other special calculation-register (like p.e. MMX).

With the content of the address-registers it's surely possible to calculate too, but their values mostly refer to values in memory, where the predefined reference-values of the are located. Therefore the initial conditions of the address-register-values should only gyrate circadian through the references.

The Status_{*n*}/EFlags_{*n*}-Register higher bytes are always filled with the same initial conditions from SAC.*actual_Processor_Mode*. The ConditionCodes in the lowest byte of Status_{*n*}/EFlags_{*n*} have got variable initial conditions. With the Control-, Debug- and MachineState-Registers and the other special registers is dealt carefully too - the always get the same default-values.

For every form initial values by opcode-execution changed (destination-) register, in a loop over all possible source-registers it's ascertained which possible operation-types of the OIT could caused the value in the changed (destination-) register.

For every appropriate source-destination-register-combination a table-entry is generated (unitary operators get the *Register_ID_source*=1) and for every matching operation-type between possible source and destination the OIT-belonging *Operations_BitCode*-bit is set [p.e. $2+2=2*2=2^2=2\ll 1=2|4=\dots=4$ for one or two source registers 2 (the other could be a constant) and one destination-register 4].

The ORT-columns are described in fig.5 and fig.19 contains the value-assignment-algorithms.

1.3.2.5 The OpCode-Learn-Table [OLT - Fig.6]:

The OLT is a subsumption of the effects of the actual opcode on the various used initial conditions. In the first 6 columns informations of fatal effects of opcode-execution are collected. Then there's the difference between the instruction-pointer=program-counter -value after and before opcode-execution and after it the condition-codes which could have caused a jump [redundant to $ICT.Register_Value(Register_ID=0)$].

Then in *Register_changed_BitCode* and *Register_source_BitCode* the Bits of all possible destination- and source-registers follow out of the belonging ORT-entries and after it in *max_Operation_BitCode* and *min_Operation_BitCode* the bitwise ORed and ANDed BitCodes of the ORT.*Operations_BitCode* -Entries.

The duration and time of opcode-execution are stored, and in *aim_valuation* analog VFT.*Valuation_Function(ADT.aim_Valuation_FunctionID)* it's appraised, how valuable the opcode-execution is in reference to reaching the programming-aim (value-assignment is shown in fig.20).

1.3.2.6 The OpCode-BaseTable [OBT - Fig.7]:

The opcode-base-table shows the ascertained total effect of one opcode in reference to all initial conditions. Fig.21 explains, how the evaluation (column-filling) happens, for so getting a "warrant of apprehension" of the opcode.

The OBT contains the resume of all executions of the opcodes, which later is necessary for the aim-directed planning of code combinations.

1.3.2.7 The Combinations-RegisterTables [CRT(i) - Fig.8]:

The Combinations-RegisterTables are created dynamically, built analog the ORT, with the difference, that effects of opcode-combinations are analysed here. So the CBT(2) has got one more opcode in the primary key than the OBT = CBT(1), and CBT(3) has got three opcodes, etc.

1.3.2.8 The Combinations-LearnTables [CLT(i) - Fig.9]:

Here the same is valid, like in the CRT(i). Analogous the OLT, one CLT(i)-row contains the effects of one opcode-combination in reference of the used initial conditions.

Here first the column *CLT(i).gradient_aim_valuation* gains importance. While it was identical to *aim_valuation* in $CLT(1)=OLT$, in $CLT(i \geq 2)$ it contains: $CLT(i).aim_valuation - CLT(i-1).aim_valuation$ (see third line of fig.20).

1.3.2.9 The Combinations-Base-Tables [CBT(i) - Fig.10]:

Analogous the CBT(i) show the resume of the effects of all initial conditions in reference to one opcode-combination. The value-assignments are shown in fig.21. CBT(n) (where n is the largest i) is the combination-plan-table - it's the place where the aim-solution-program originates.

If $ADT.aim_fulfilled_Flag_Function(CPT-PK)=1$ (TRUE), the solution-program of the given aim is found and it'll be enrolled into the AST.

1.3.2.10 The Aim Solution Table [AST - Fig.11]:

For every given programming abandonment the found solution-programs, their lengths, execution-times and used registers and operations (→bitcodes) are enrolled here.

1.3.2.11 The Aim Description Table [ADT - Fig.12]:

The ADT assigns to every programming aim an identification-number, a short description, one bitcode-combination of the source- and one of the destination-registers which should be used (if possible) and one bitcode-combination of forbidden source- and one of forbidden destination-registers; further a string of former solution-programs which could be implemented, and a aim-solution-flagfunction which returns TRUE if the opcode-combination (program) solves the problem for the desired source- and destination-registers and finally an identifier which references to the valuation_function in the VFT, that appraises the closeness to the complying programming-aim, which is among others dependant of this aim-solution-flagfunction.

1.3.2.12 The Function-Identification-Table [FIT - Fig.13,14]:

In the FIT basic subfunctions are provided, which can be used for composing the energy valuation-function.

It's introduced in two variations:

- a.) for generating a dynamical valuation function in SQL,
- b.) for generating a dynamical valuation function in machine-code.

The alterable building of a valuation function is easier to accomplish in SQL, but the execution-time in machine-code is much faster and every new composed SQL-valuation-function has to be parsed again.

In the future the valuation-function should only be composed in machine-code. This has the additional advantage that the AC-program could use some solved solution-functions again as subfunctions in the FIT for later use for composing the valuation-function.

1.3.2.13 The Valuation-Function-Table [VFT - Fig.15]:

The VFT contains the dynamic valuation-system in reference to the own "well being" (energy-register) and to the closeness to the programming aim(s).

The VFT.Valuation_Function(Type='E', SAC.Energy_Valuation_Function_ID) appraises energyspecific actions and the Valuation_Function(Type='A', SAC.Aim_Valuation_FunctionID) the closeness to the programming aim(s).

The VFT.Function_ID_Chain contains the concatenation of the FIT.Function_ID's, that means the execution-chain of the subfunctions: Here causes NUM (see fig.13b), that the following value is used as a number of byte-length, VALUE denotes that the following value is the column-number of the CPT=CLT(n), from which actual row the value is taken, EREG means the Register_ID of the energy-register, S/D_REG denotes the value out of ADT.all_source/dest_Registers_BitCode and AIM_F is the

result of the ADT.aim_fulfilled_Flag_Function. The unitary operations operate on the last result of the Function_ID_Chain and the binary operations on the last two results.

On every accommodation, enhancement, or other amelioration of these valuation-functions the Valuation_Function_ID is incremented and a new entry with the modified Valuation_Function is created and the efficiency of all valuation-functions are reappraised:

$VFT.Valuation_Function_value = SAC.Energy/Aim_self_valuation_Func(...)$, for to have an efficiency-gradient for further improvements.

The functionality of the dynamic valuation-system is described in 1.3.7.

1.3.2.14 The Status of the AC-Program [SAC - Fig.16]:

This table has no primary key and only one row. It contains status-informations of the AC-Program and two self-valuation-functions, which appraise the efficiency of the energy-valuation-function and of the valuation-function of the programming-aim-closeness (VFT) by evaluating the range of their valuation-results.

These self-valuation-functions are, in opposite to the energy- and aim-closeness valuation-functions, not modifiable by the AC-program itself, but can be changed by the user.

1.3.2.15 The Energy-Learn-Table [ELT - Fig. 17]:

In the ELT data are stored about all energy relevant actions over the actual initial conditions, that means for all opcodes and code-combinations, which pertain the last data-register.

The valuability of an energyspecific action is appraised according to $ELT.Energy_valuation = VFT.Valuation_Function(SAC.Energy_Valuation_Func_ID)$.

1.3.2.16 The Energy-Base-Table [EBT - Fig.18]:

Like in the CBT(i) for programming-aim-closeness, in the EBT the effects of energy-changing opcode-combinations for all initial conditions are collected.

1.3.3 Preparing the initial State of the System

For to reset the system later into the initial state without booting, several pointers have to be latched. Afterwards all exception-vectors are intercepted by own routines, because of the initial trying to use arbitrary numbers as machine-opcodes, although many of these trying cause fatal exceptions, because they're illegal opcodes (not usable) or the opcode causes an exception on one of the initial conditions. Abnormal system end would be the consequence, if not all exception-vectors would be captured.

In the case that the AC-program should run preclusively, you'll have to

a.) stop multitasking by disabling it by an operating-system routine or by setting the IRQ-mask of

the processor to NMI.

- b.) save all system-exception-vectors.
- c.) set all system-exception-vectors to own analysing and handling routines.
or if it should later run with other programs or perhaps with further AC-programs:
 - a') increase own task-priority.
 - b') save all task-exception-vectors.
 - c') set the task-exception-vectors of the AC-program to own analysing and handling routines.

- d.) save the statusregister_μ ($\triangle EFlags_{\mu}$) and the user-stackpointer.
- e.) save the values of the other address-registers and of the data-registers.
- f.) save the values of the segment-, control-, debug- and special-registers.
- g.) set exception-vectors, which load additional data to the supervisor-stack (p.e. on address-violation-exception several processors load additional information like access-address and opcode onto the supervisor-stack) to a this fact considering interceptor-routine.
- h.) set the privilege-violation-exception-vector to a special capturing routine.
- i.) set one trap-vector to a routine, where the system should continue inside the supervisor-mode after this trap occurs.
- j.) execute this trap intentionally to change the CPU-mode from user-mode to supervisor-mode (the system now continues on the above set routine).
- k.) set the trace-exception-vector to an own trace-routine for later effect-analysis after number-as-opcode execution.
- l.) set the bits in the first word on the supervisor-stack so, that when the SR is loaded from SSP, the trace-Flag is set and the IRQ-mask is set to NMI (p.e. this is #8700 on Motorola) because while the following base-opcode-learning no interrupt should be possible and after execution the effects should be analysed.

See referring to this fig.24a.

1.3.4 Base-learning from execution of all single opcodes:

1.3.4.1 OpCode generation and execution:

- a.) Generate a 32-Bit-Number as an opcode, starting on #0000.0000, later increase by 1. [if you already know the CPU instruction set, you can skip the opcodes which would overwrite memory (p.e. memory-move-commands).]
- b.) Set data- and address-registers, and the address-register destination-values and the values one DWord below on predefined test-values (initial conditions) and clear the condition-codes in $EFlags_{\mu}$ (or use several initial conditions too).
- c.) Write the above generated opcode into the test-location in memory. Then fill the memory behind with zeros up to the maximum possible opcode-length, if zeros mean the mnemonic

"ORI #0, Reg.0" (or another effectless command) or fill with NOPs.

This is necessary because on a long command the zeros are not meaningless [and then often less destroying (p.e. memory overwriting) than the NOP-corresponding opcode-number], and on a few processors a clearing of the trace-flag is possible in the while execution used user-mode too, which effects the execution of the following numbers as opcodes too [if now the zeros are not effectless, NOPs have to be used to prevent further effects (like memory- or program-selfoverwriting)].

After these zeros or NOPs the Trace-Bit-Cleared handler is following.

d.) Set content of the supervisor-stack, which is loaded by return from supervisor-mode into important user-registers like EFlags_u/Status-Register_u, IP_u/PC_u, etc., so, that CCR will be cleared or set on initial conditions, IRQ mask will be set to NMI, the trace-bit will be set and the supervisor-bit[mask] will be cleared and the DWord behind is the location of the test-opcode.

Now execute the return from supervisor-mode by the corresponding opcode-command (p.e. RTE_d): EFlags_u/Status-Register_u is now loaded by above described values and the test-opcode executes, because IP_u/PC_u is loaded by its address.

- If now a fatal exception occurs (except trace, p.e. privilege violation, etc.), the kind of exception is briefly shown graphically if desired, and it will be continued by generating and executing the next opcode.
- If an initial condition dependant exception occurs (like address error, division by zero, ...), a relation between initial condition an exception is analysed [attention: on several exceptions, because of trace, an in many literature not documented combination of both exceptions (internal processor handling) can occur (p.e. using M68000 on Trap, Chk, Div/0 in combination with Trace).]
- If no exception occurs (not even trace) the opcode-execution cleared the trace-bit (should never occur while executing single opcodes) and the handler behind the opcode is executed.
- On Trace-Exception (normal case) a usable opcode was generated which execution-effects have to be analysed now.

1.3.4.2 Analysis of opcode-repercussion and saving the analysis-result:

- a.) After execution the EFlags_u/Status-Register_u and the data- and address-registers and the reference-values of the address-registers an the reference-values one DWord below an the user-stackpointer are saved for analysis.
- b.) Verifying the own machine-code-checksum (of AC-Prg.) and the inactive copy in RAM (both without test-opcode placement): If checksum changed, the AC-program injured itself while executing the test-opcode (overwrote own parts of program). The corresponding corrupt-flag is set in the table. If the active version checksum changed jump into the inactive version, then compare both versions byte by byte and repair the corrupt version by replacing the bytes in the version with the changed checksum by the bytes from the version with the correct checksum.
- c.) Check the supervisor-bitmask of the saved user-stackpointer on the supervisor-stack: If the

By analysing the supervisor-bit(mask) on the supervisor-stack, it's now detectable, that before trace another low-priority-exception occurred; and by comparison of the second saved program-counter below on the supervisor-stack with the low-priority exception-vectors, now the primary exception before trace is detectable, which corresponding exception-number is stored.

- If the IP_{Ψ}/PC_{μ} increased by 5 or more bytes and less than the longest possible opcode, it was a long opcode or a short forward jump. If no registers changed from initial conditions it was a short jump.

e.) Comparison of the EFlags_{PI}Status-Register_μ and of all register-values and of the address-register destination-values and of the destination-values one max. address-length below [because of -(Adr.Reg.)], with the original-values.

In a bitmask now it's flagged which registers or its destination-values changed and it's analysed, which operations on which source- and destination-registers could have happened (heeding changes of $EFlags_{SR_{\mu}}$) and the result is stored in the ORT and OLT (see figs.5,19; 6,20) and the OBT is actualised (fig.7,21).

- f.) If it was a jump-command, in the EFlags/SR, it's analysed if it was an conditional jump.

1.3.5.1 Realisation of artificial pain:

The AC program is loaded twice into RAM. If the AC program (or another one) executes a

command which overwrites a part of the active or inactive AC program, which means an injuring of the active or inactive code (DNA), it has the ability to recognise the damage by comparing the checksum, and has now to take time to repair the damaged code by comparing damaged and undamaged code to have information about the damage-location (valuable information for test-opcode analysis) and then copying the code from the undamaged version into the injured version to heal itself. If the active code was the injured one (had the corrupt checksum), it first has to jump into the inactive duplicate to prevent errors on self-healing, because the healing-routine itself could be damaged.

1.3.5.2 Realisation of artificial hunger:

Hunger means imminent loss of energy. Energy is engendered in the cells by transforming adenosintriphosphat to adenosindiphosphat. The energy for building up adenosintriphosphat from adenosindiphosphat is gained by combustion of glucose. Missing energy (ATP) makes metabolism and with it every action, reaction on pain, or self healing on injury, impossible.

The "energy quantity" of the AC-program is modelable by the height of a value in a data register. Now it would be possible to realise hunger by decreasing the electric current to the processor by external reading of this data register and increasing an ohm-resistance inverse proportional to the register value.

A less authentional hardware-unbound solution is possible too:

Less energy is harmful for learn-process. Frugal values in the energyspecific data-register cause lower functionality on learning from opcode-executions. On less values no learning from opcode-execution is possible. And lowest values cause loss of the saved knowledge from earlier opcode-executions. If the value is zero, additional pain, which means own code injury, occurs.

On hunger the AC-program consequently has to find and execute opcodes, which increase the value of the energyspecific data-register.

Decreasing energy, which means the origin of hunger, is simulated by decreasing the energyspecific data-register by 1 after every action (=opcode-execution) [p.e. by the AC-program itself].

1.3.6 Planning on the criterions of the valuation-system:

If the system tested all possible opcodes and analysed the effects of the suitable commands, it has now the possibility to learn planning aimed to comply its basic needs or its programming-aims: Therefore it combines the opcodes, executes them using all initial conditions and analyses, what effected the code-combination on every single initial condition.

Because mostly longer opcode-combinations are necessary to fulfill the programming-aim, it plans the code combination by using only codes which caused no injury (own damage) or better no RAM access at all and caused no fatal exceptions (divide-error or overflow-exception are allowed) and

used no forbidden registers or opcodes (ADT.*unused_Registers_BitCode*| ADT.*unused_Operations_BitCode*). OpCodes which use the desired destination and source-registers are preferred (ADT.*all_source/dest_Registers_BitCode*).

ADT.*aim_fulfill_valuation_mode* appoints, if the valuation-function is existent in SQL or directly in machine-code. For the beginning user the slower SQL-version is more convenient and the specialist would prefer to use an assembler-*aim_fulfilled_Flag_Function* (ADT) which is transformed into machine-code, because it's much faster on complex valuation-functions than SQL.

1.3.7 The dynamic-reflexive valuation-system:

1.3.7.1 Valuation of the programming-aim closeness:

The ADT.*aim_fulfilled_Flag_Function*(Aim_ID), returns TRUE, if the programming-aim is obtained and the VFT.*Valuation_Function*(Type='A', ADT.*aim_fulfilled_Flag_Function*, VFT.*Function_ID_Chain*), supplies a signed-byte value, which means the closeness of the actual CLT(n)-opcode-combination on the used initial conditions to the aim-solution. The result is stored into CLT(n).*aim_valuation* and builds up in comparison with the last CLT(n-1).*aim_valuation* the gradient CLT(n).*gradient_aim_valuation*.

Because of the solution program has to work for all initial conditions, the maximum- and the average valuability of the opcode-combination, both as average over all initial conditions, is stored in CBT(n).*max_aim_valuation* and CBT(n).*avg_aim_valuation* ; and the gradients to the corresponding values of the last CBT(n-1) build up CBT(n).*max_grad_aim_valuation* and CBT(n).*avg_grad_aim_valuation*.

If the boundary-value of -128 or +127 was the valuation-result, the VFT.*boundary_value_counter* is incremented, and analog *low_value_counter* is incremented, if a valuation-result between -16 and +15 occurred.

Using these statistical data, and on an analysis of all CLT(i).*aim_valuation* -values, p.e. if you count the number of values in a small value-range running from min-possible-result to max-possible-result (-127 to +128) in dependence of the width of the value-range-window, the SAC.*Aim_Self_Valuation_Func* appraises after every programming-aim attainment the valuation-results of the VFT.*Valuation_Function* and with its efficiency.

If the most valuation-results of the VFT.*Valuation_Function* p.e. were near the boundary-values (min./max.), the valuation-function was too steep and has to be flattened, which means inside the VFT.*Function_ID_Chain* have to be more elements with negative FIT.*Function_Flatten*. The opposite is valid, if the most valuation-results caused a high VFT.*low_value_counter*.

So after every solution of a programming-aim a self-valuation of the valuation-function occurs and a further step in the self-programming of the valuation-function. New elements are added to the valuation-function and sometimes elements are omitted and the steepness is adapted.

Then the valuation is done again and it's revised if the new valuation-function would have supplied

a better range of valuation-results.

If the new range of valuation-results was worse than the one before (valued by *SAC.Aim_Self_Valuation_Function*) then the modification of the valuation-function is quashed and another modification is tried. If the changing of the valuation-function ameliorated the range of valuation-results and the self-valuation-function returns a positive value, then it's continued with the next programming-task - otherwise a further amelioration of the valuation-function has to be done until self-valuation returns a positive value.

1.3.7.2 The dynamic energy-valuation-function:

The dynamic energyspecific valuation-system runs as follows:

0.) Because of the results of the energyspecific valuation-system are constricted to the range of *signed_byte* the valuation-function is embedded into a frame:

valuation-result := MIN[MAX(valuation-function, -128), +127]

1.) The energy-valuation-function of 0-th order is "how saturated am I after the action ?":

valuation-function(0) := MIN[MAX(Energy_after, -128), +127]

2.) The valuation-function of 1st order is "how much more saturated am I after the action than before ?":

valuation-function(1) := MIN[MAX(Energy_after - Energy_before, -128), +127]

3.) Because of the energy-register is of the type *unsigned integer* (DWord), the boundaries of the valuation would too often the result. Therefore a kind of logarithm or ...

valuation-function(2) := MIN[MAX[SQRT(Energy_after - Energy_before), -128], +127]

4.) Now negative energy-gradients would cause wrong signs, therefore 3rd square or ...

valuation-function(3) := MIN[MAX[SGN(EnergyGrad) * SQRT(EnergyGrad), -128], +127],
where EnergyGrad = Energy_after - Energy_before

Possibly the function $\frac{1}{2} \cdot \text{SGN}(\text{EnergyGrad}) \cdot \text{SQRT}(\text{SQRT}(\text{EnergyGrad}))$ would be better, because it reaches exactly to the boundary-values, but maybe the boundary-values are reached very seldom and a subtler structuring around zero would be more important.

This depends how often the boundaries are reached and how much energy-gradients supply small values. Maybe the naked result-value (*Energy_after*) has to be weighted stronger and a valuation of the gradient alone is not sufficient. Moreover it had to be considered how much and which further registers are concerned beneath the energy-register and which and how much types of operations were executed, etc., and finally the execution-time of the energy-valuation-function itself. Therefore the energyspecific valuation-system has to be rarefied and adapted (like the valuation-system of intelligent biological life forms).

Using dynamic embedded [PL]/SQL the changing and reparsing of the as string stored valuation-function no problem. Because of the execution-speed and the possibility of implementation of earlier programming-aim solutions the energy-valuation-function should be used in machine-code prospectively.

The task of the amelioration of the energyspecific valuation-function is done, like the programming-

aim specific one, after every fulfilling of a programming-aim.

Valuation-system and valuation-results are always reflexively.

1.3.8 Reaching Self-Consciousness, Reproduction and Evolution:

Through the process of self-healing on pain/damage the program knows its location in memory. It now has the possibility to check out the effects of its own opcodes one after another, then consecutive opcode-combinations etc., and when it finally knows the effect of its total length, it'll reach self-consciousness and then has the possibility to reproduce itself and to ameliorate itself awarely while reproduction using its acquired knowledge (p.e. by removing the incommodious decrementing of the energy-register).

The intelligent conscious-varying reproduction is much more preeminenced than the biologic-genetic one, because the latter only refers to existing genes/DNA, while the AC-program can vary itself by changing and extending its code intentionally using its "experience of life".

1.4. Conceptual Formulation of the Programming-Aim and Examples of Achievement

To the AC-program an arbitrary programming-aim is challenged by giving one or more criterions in *ADT.aim_fulfilled_Flag_Function*, by which it can check out, if it solved the task.

Its job is it now to develop a program, which solves the problem for all initial-conditions.

1.4.1 Example_1: Developing a Program to compute the average:

A very simple but easy to comprehend job for the AC-program could be: "write a program that computes the average over two integer-variables".

The AC-program fulfills this task, if the difference between the result and the lower number is equal to the difference of the higher number and the result, and this is functioning for any arbitrary input-numbers.

But the AC program doesn't know the instruction-set of the processor - it knows now only the opcodes which caused no damage and no fatal exception and it knows the opcode-effects in reference to the different initial-conditions.

Through corruption-selfhealing or the energy-register it already knows easiest abandonments like "execute an action which causes no pain" or "execute an action which makes me saturated".

For attaining economic programming-aims, it now needs valuation-variables, which reveals answers to the following questions:

- a.) How much nearer or farther away from the programming-aim took me the last added opcode-combination (that every added single opcode of it may cause opposite effects is irrelevant).
- b.) How many processor clock-cycles needed the solution-program.
- c.) How many bytes long is my program and how many opcodes are included ?

These answering variables (CBT-table-columns) are:

aim_valuation ; cycles_of_execution ; OpCode_length_or_jump.

The input-variables in the example-task could be in the first two data-registers (EAX_π, EBX_π respectively D0_μ, D1_μ respectively GPR0_ψ, GPR1_ψ), here *R0* and *R1*.

The return-value should be the third data-register (ECX_π|D2_μ|GPR2_ψ), here *R2*.

If the task is solved for arbitrary input-values, the program is ready, because it's function.

If there are more than one solution, the one is chosen which needs less clock-cycles.

The task-specific aim-fulfilled valuation-function, which computes OLT.*aim_valuation* is consequently in this example:

ADT.*aim_fulfilled_Flag_Function*(average of R0 and R1) = { (R2-R0)=(R1-R2) }

Here the problem could occur, that one input-value is even and the other one is odd, which means that this input-combination would have no solution. To implement that, the *aim_fulfilled_Flag_Function* should be enhanced for integer-variables: { (R2-R0)=(R1-R2) || (R2-R0)+1=(R1-R2) }.

The AC-program will find several solution-programs and takes the one with the least needed clock-cycles.

A possible solution would be in CBT(3): MOV R0,R2 ; ADD R1,R2 ; SHR R2

(... naturally in machine-code of the used processor - using a Pentium that would be the 48-bit-number #89C2.01CA.D1EA, using a Motorola that would be #2400.D282.E2C2 and using a PowerPC (RISC) a 96-bit-number would be the solution).

1.4.2 Example_2: generation of a programs for computation of the cube-root:

A further easy development-task would be "write a program that returns the cube-root of a FFP (fast floating point) number"; the input-variable should be R0 (EAX_π) and the output-variable R3 (EBX_π).

The AC-program fulfills this task, if the result multiplied with its square is equal to the input-value (and this is valid for all initial conditions):

⇒ *aim_fulfilled_Flag_Function*(cube root) = { (R3*R3*R3)=R0 } (←naturally in FFP-multiplier.)

A single command like the square-root (FSQRT) doesn't exist for the cube root.

The solution-program could be in CBT(8) using a Pentium II as follows:

Op1(16b):	MOV CL,3	;ECX=\$????:0003	[1011.0001:0000.0011]
Op2(16b):	FLD1	;ST(0)=1.0	[1101.1001:1110.1000]

Op3(16b): FIDIV CX	;ST(0) = $\frac{1}{3}$	[1101.1110:1111.0001]
Op4(16b): FLD EAX	;ST(0) = RO ;ST(1) = $\frac{1}{3}$	[1101.1001:1100.0000]
Op5(16b): FYL2X	;ST(0) = $\frac{1}{3} \log_2(RO)$	[1101.1001:1111.0001]
Op6(16b): FLD1	;ST(0) = 1.0 ;ST(1) = $\frac{1}{3} \log_2(RO)$	[1101.1001:1110.1000]
Op7(16b): FSCALE	;ST(0) = $1.0 * 2^{\lceil \frac{1}{3} \log_2(RO) \rceil}$	[1101.1001:1111.1101]
Op8(16b): FST EBX	;EBX = $2^{\lceil \frac{1}{3} \log_2(RO) \rceil}$	[1101.1001:1101.1011]

(... naturally only the second of these columns as a 128-bit-number with the bits set as shown in the last column.)

Hexadecimal that would be: B103.D9E8.DEF1.D9C0:D9F1.D9E8.D9FD.D9DB.

This would be a possible solution-number (=program) for the given task (there're surely shorter and faster solutions too).

1.5. Needed Hard-disk-Space and Oblivion

In both examples 16-bit-opcodes would be sufficient, but it's obvious, that large programming-aims would need much hard-disk space. Therefore the AC-program has to forget unimportant or error-causing opcode-combinations.

1.5.1 Table sizes:

IST, RIT and CIT need neglectable disk-space.

Theoretically there could be $size(OBT) = 2^{32} * \sum \text{bytes}(\text{column}(i)) = 485 \text{ GB}$, but also on a RISC processor never all 32-bit-combinations are used as a valid opcode and realistic are as an average on RISC processors about 28 bits $\Rightarrow 30 \text{ GB}$ and on CISC processors about 20 bits $\Rightarrow 118 \text{ MB}$ (on latter the most are 16 bit-opcodes, there're a few 8-Bit- and several 24- and 32-bit-opcodes, and the ones which are longer 32-bits are not used (we don't need memory-to-memory-operations for example and the functionality of one long opcode can be substituted by two or more shorter opcodes)).

The 62 initial conditions can cause $size(OLT) = 2^{[20..28]} * 62 * \sum \text{bytes}(\text{column}(i)) = 3 \text{ GB}$ (CISC) up to 832 GB (RISC) and $size(ORT)$ could reach the same quantity, considering that one opcode mostly pertains one destination- and one source-register (unitary ones use only the destination-register and a few seldom opcodes use 3 or more registers). But there could be many effect-belonging *Operation_BitCodes*, which would increase the tablespace dramatically, if it wouldn't be compensated by the many opcodes which cause an exception, where less information has to be stored.

A much larger problem is the exponential growth of the $CxT(i)$, because every i multiplies the needed tablespace by factor of $2^{[20..28]}$. But this is exact that what should be compensated by the dynamic valuation-system. It decreases its knowledge-absorption tolerance referable to the

remaining hard-disk space: So opcode-combinations with least $CBT(i).max_aim_valuation$ or $.avg_aim_valuation$ are forgotten and so will not be combined with other opcodes.

Although if the demand of hard-disk space arrears large, this is no problem in near future. Also the according to the combination-possibilities and table-sizes increasing calculation-times are compensated by larger and faster becoming hard-disks and the increasing efficiency of the processors.

1.5.2 Oblivion:

Like all intelligent life-forms, the system has to forget unimportant and less important informations, because

- a.) the disk space is limited, and
- b.) data access time becomes slow on very large data-tables.

Therefore after every satisfactorily achievement of objectives, when a new programming-aim is given, the ELT and all $CxT(i)$ -tables over a remaining disk-space dependant i are deleted, and the opcode-combinations in the remaining $CxT(i)$ are revalued referable to the new programming-aim and forgotten opcode-combinations are added, if they're valuable for the new aim and the higher $CxT(i)$ are recreated dynamically.

1.6. Becoming Conscious

Through `try_and_error` the program learned what effectuated every action and what are the effects of which sequence of actions.

Through corruption-healing (if own code was overwritten in RAM) it had to repair its code and so it knows its position in memory.

If it once knows the effect of its own machine-code, it gets self-consciousness and has the ability to reproduce its code and to ameliorate it while reproduction.

Through the so initialised evolution the AC becomes complexer and better and will be able to solve larger and larger programming-tasks in future.

1.7. Presentment of the Economic Advantages

Here a totally new area of using a computer is presented. While normally in a computer run by man generated programs, which execute user-controlled applications, the AC-program itself develops and executes programming-aim oriented routines, which can later embedded into a large application.

The demand for software-development is worldwide much larger than the human potential of developers.

A system which learns to write programs itself has the capability to solve smaller development-tasks, and will in future, after several evolution-steps, have the capability to develop complex programs as the solution of large tasks too, if it has got enough hard-disk space.

The programmers will not have to develop all the routines they need - they order the routine from the AC-program and embed it into their application. So the companies can finish and sell their software-products earlier.

CLAIMS:

What is claimed is:

1. A method using a computer which automatically generates and executes machine-code, comprising the steps of
 - a) preventing the multi-tasking of the operation-system, by setting the interrupt-mask of the processor to NMI or using a multitasking-disable-routine of the operating-system;
 - b) capturing the processors exception-vectors by own analysis-routines;
 - c) generating normal numbers and writing them into memory;
 - d) backing up the current values of the processors registers;
 - e) positioning the instruction-pointer=program-counter to the generated number in memory and executing the number like it would be a processor-opcode; and
 - f) analysing the effects of this opcode-like execution of the number and storing the analysis-results, p.e. in a database.
2. A method according to claim 1, wherein said step of (d) "backing up the current values of the processors registers" comprising the steps of:
 - a) not only saving them for a later comparison but setting them to predefined initial conditions;
 - b) setting them not only one time but several times to many different predefined initial conditions which means several executions of one number in step (e) by these different initial conditions for to have a more efficient analysis-determinations of possible number-execution-effects in step (f).
3. A method according to claim 2, wherein said step (1f) "analysing the effects of this opcode-like execution of the number" further comprising the step of determining the number-opcode's mnemonic and its related source- and destination-registers by regarding all execution-effects of every initial condition.
4. A method according to claim 3, wherein said step (1c) "generating normal numbers and writing them into memory" comprising the steps of combinations, which means:

- a) taking one number which analysed results are already stored and appending another number with stored analysis-results to analyze the execution-effects of this two-number-combination and store the result.
 - b) combining 2-number-combinations, which effects are already analysed, with a further analysed number and analysing and storing the analysis-results of the effects of these 3-number-combinations.
 - c) combining a 3-number-combination, which effects are already analysed, with a further number, which effects are already analysed, or combining two 2-number-combinations, which effects are already analysed, and analysing and storing the effects of these 4-number-combinations.
 - d) combining larger combinations, which effects are already analysed, with numbers or combinations, which effects are already analysed, and analysing and storing the effects of these larger combinations.
5. A method according to claim 4, further comprising the step of using only combinations for further use, which got a positive value from a valuation-function, which appraises the valuability of the combination in reference to reaching a pregiven programming-aim, not causing fatal exceptions, not overwriting exception-vectors or the program, avoiding to use forbidden registers or extensive writes to memory or large jumps, etc.
 6. A method according to claim 5, further comprising the step of valuating and changing the valuation-function of the dynamic valuation-system by a meta-valuation-function valuating the results of the valuation-function according to clustering to boundary-values, low-values, other fixed values, etc., and then revaluating the results of the new valuation-function.
 7. A method according to claim 4, further comprising the step of implementing calls to operation-system routines which are offered in a table with entrance-address and source- and destination-registers.
 8. A method according to claim 7, further comprising the step of using only calls and combinations for further use, which got a positive value from a valuation-function, which appraises the valuability of the call-combination in reference to reaching a pregiven programming-aim, not causing fatal exceptions, not overwriting exception-vectors or the program, avoiding to use forbidden registers or extensive writes to memory or large jumps, etc.
 9. A method according to claim 8, further comprising the step of offering the disassembly of the solution-programs which solved the programming-aim.
 10. A method using a computer which automatically generates and executes machine-code,

comprising the steps of

- a) capturing the tasks=processes exception-vectors by own analysis-routines;
- b) generating normal numbers and writing them into memory;
- c) backing up the current values of the processors registers;
- d) positioning the instruction-pointer=program-counter to the generated number in memory and executing the number like it would be a processor-opcode; and
- e) analysing the effects of this opcode-like execution of the number and storing the analysis-results, p.e. in a database.

11. A method according to claim 10, further comprising the steps of modelling the following basic needs:

- a) "no_pain", where pain means damage to the own program, which is an overwriting of the own machine-code, which is recognised by comparing the programs checksum after every execution of a number or number-combination, and repairing damaged parts of the own machine-code from a duplication, which causes a decrementation of the energy-register for every damaged opcode of the own program which now has to be copied for reparation; and
- b) "no_hunger", where hunger means the imminent loss of energy, where energy is modelled by the value of a predefined register, which causes negative effects on low values like
 - the loss of the capability of appraising combinations referable to the programming aim on low values of the energy-register,
 - mistakes on the valuation of the combination-execution concerning the source-registers on very low values of the energy-register,
 - the loss of the capability of self-repairing on "pain" on extreme low values of the energy-register,
 - a hardware-dependant decreasing of the power supply of the RAM (p.e. by increasing a resistor) on two times in series extreme low values of the energy-register,
 - a hardware-dependant decreasing of the power supply of the processor (p.e. by increasing a resistor) on three times in series extreme low values of the energy-register,and this energy-register is decremented after every action, where action means the execution of a number.

12. A method according to claim 10, wherein said step of (10c) "backing up the current values of

the processors registers" further comprising the steps of

- a) not only saving them for a later comparison but setting them to predefined initial conditions;
- b) setting them not only one time but several times to many different predefined initial conditions which means several executions of one number in step (10d) by these different initial conditions for to have a more efficient analysis-determinations of possible number-execution-effects in step (10e).

13. A method according to claim 12, wherein said step (10e) "analysing the effects of this opcode-like execution of the number" further comprising the step of determining the number=opcode's mnemonic and its related source- and destination-registers by regarding all execution-effects of every initial condition.
14. A method according to claim 13, further comprising the step of valuating the effects of the execution of the number in reference to its basic needs, which means positive values for numbers, which cause no pain or increase the energy-register and negative values for numbers which cause above defined "pain" or "hunger".
15. A method according to claim 14, further comprising the step of combining numbers and executing these combinations and valuating the effects of the executions of these combinations.
16. A method according to claim 15, comprising the step of running several of these said programs as an extra process=task.
17. A method according to claim 15, comprising the step of using a network topology where on two or more of the networked computers is running one of these programs.
18. A method according to claim 14, further comprising the step of analysing a pregiven code by executing larger getting sequences of it instead of executing number-combinations, for to evaluate the effects these code-sequences or at least of the complete program.
19. A method according to claim 18, further comprising the step of improving the analysed code in the direction of a pregiven programming-aim.
20. A method according to claim 15, further comprising the step of implementing calls to operation-system routines which are offered in a table with entrance-address and source- and destination-registers.

ABSTRACT OF THE DISCLOSURE:

A method for generating a simple kind of computer based artificial consciousness, which means to give a in a computer running invention-pursuant program the capability to act and to know the effects of its actions and to plan further actions consciously. This is realized by giving the computer the capability to program its processor by its own and to plan that self-programming targeted. This works, because the computer learns to program in machine-code by its own and it has got a dynamical valuation system to weight if its actions are useful or not. Further basic needs like "no_pain" and "no_hunger" are modelled to make it act to fulfill its basic needs. It has also the capability to solve a pregiven by several formulas determined programming aim, which means to develop logical programs which then can be implemented by users into their projects.

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DRAWINGS:

3.1 Relational Database of the AC-knowledge:

3.1.1 ER-Diagram of the AC-Database:

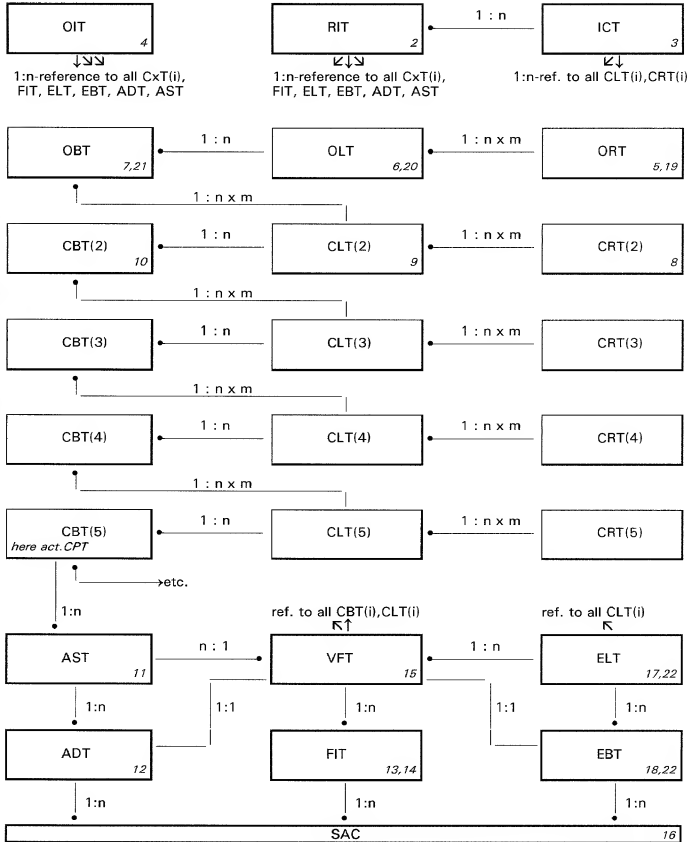


Fig. 1

3.1.2 Tables of the AC-Database:

Register-Identification-Table: [RIT: every processor-register gets a Reg.ID and a Reg.BitCode]
column: datatype value-range meaning:

Register_ID (PK)	signed byte	-128..127	0 = Flags-reg., 1... = data-reg., adr.reg., adr.reg.-destinations, FFP-reg., control-reg., debug-reg., etc.; neg.reg.ID = exception-vector-nr. (no processor-reg.)
Register BitCode	number	0..2 ¹²⁸⁻¹	2*Register ID, 0 if Register ID is negative.
Register_type	char(1)	1 Byte	type of register: # = flags-reg., D = data-reg., A = address-reg., V = adr.-reg.-destination, E = exception-vector, etc. (processor dependant).
Register number	byte	0..127	current number of the {register-type}-registers.
Register_Description	varchar2(32)	≤32 Bytes	optional description of the register[reference].

Fig.2a

Register ID	Register BitCode	Register type	Register number	Register Description
...	0	E	...	{for all Exception-Vectors}
-8	0	E	8	Privilege-Violation Exception
...	0	E	...	{for all Exception-Vectors}
0	1	#	0	Status-Register (Δ EFlags ₈)
1	2	D	0	Data-Register D0
...	4	D	...	{for all Data-Registers D1-D6}
8	8	D	7	Data-Register D7
9	16	A	0	Address-Register A0
...	...	A	...	{for all Address-Registers A1-A6}
16	65536	A	7	Address-Register A7 (= USP)
17	131072	V	0	Destination of Address-Register A0
...	...	V	...	{for all Adr.-Reg.-Destinations A1-A6}
24	16777216	V	7	Destination of Adr.-Reg. A7 [= (USP)]
25	33554432	v	0	Destination before Adr.Reg A0 [- (A0)]
...	...	v	...	{for all Adr.-Reg.-Dest. before A1-A6}
32	\$1.0000.0000	v	7	Destination before Adr.Reg A7 [- (USP)]
33	\$2.0000.0000	F	0	Floating-Point Data-Register FPRO
...	{for all further Registers}

Fig.2b (RIT in a Motorola-Example)

Initial-Conditions-Table: [ICT: every register-(reference) gets 62 initial conditions]

column:	datatype	value-range	meaning:
IniConNr (PK)	signed byte	-31..+30	number of the initial-condition combination
Register ID (PK)	signed byte	0..127	see RIT
Register Value	integer	0..2 ³²⁻¹	value of the register-(reference) on actual IniConNr

Fig.3a

Therefore 62*#registers test-values are generated, p.e. using the following function:

$$\text{Register_Value}(\text{IniConNr}, \text{Register_ID}) = \text{SGN}(\text{IniConNr}) * \text{INT}(\text{2} * \text{ABS}(\text{IniConNr}/2) + \frac{1}{2}) + \text{prime_number}(3 * \text{Register_ID}) \quad \text{or similar.}$$
 ,where prime number(0) = 0 and prime number(-n) = -prime number(n)
 [no 2 equal register-(destination)-values]

Fig.3b

Operation-Identification-Table: [OIT: every processor-operation gets an Oper.ID and an Oper.BitCode]
 column: datatype value-range meaning:

Operation ID (PK)	signed byte	-1..63	bit of the Calculation BitCode - see table below.
Operation BitCode (FK)	number	0..2 ¹²⁸ -1	2 ⁿ Calculation ID - see table below.
Operation Type	char(5)	5 Bytes	5 characters-code of the operation-type, see Fig.4c
Operation Mnemonic	char(5)	5 Bytes	abbreviation of the operation - see table below.
Operation Description	varchar2(32)	≤32 Bytes	optional description of the operation - see table below

Fig.4a

Operation ID	Operation BitCode	Operation Type	Op.Mnemonic	Operation Description
-1	0	???	unknown operation
0	1	.I11?	TST	set flags in dependence of reg.(-ref.)
1	2	.I12!	NEG	negation amount
2	4	.I12!	NOT	bitwise inversion
3	8	:I02	MOVI	const.integer→register(reference)
4	16	:I12+	ADDI	add constant integer
5	32	:I12-	SUBI	subtract constant integer
6	64	:I13*	MULI	multiply constant integer
7	128	:I23/	DIVI	divide by constant integer
8	256	:I13%	MODI	rest of integer-division
9	512	:I12*	SHLI	integer-times a duplication
10	1.024	:I12/	SHRI	integer-times a halvation
11	2.048	:I12	ORI	set bits set in a constant integer
12	4.096	:I12&	ANDI	clear bits not set in a const. integer
13	8.192	:I12?	BTSTI	check if int-th bit is set in reg.(-ref.)
14	16.384	:I12?	CMPI	reg.(-ref.)-comparison with integer
15	32.768	II22	MOV	move src.-reg.(ref.)→dest.reg(ref.)
16	65.536	II22+	ADD	addition of register(reference)
17	131.072	II22-	SUB	subtraction of register(reference)
18	262.144	II23*	MUL	multiplication of register(reference)
19	524.288	II33/	DIV	division by register(reference)
20	1.048.576	II33%	MOD	rest of division by register(ref.)
21	2.097.152	II22*	SHL	register(-ref.)-times a duplication
22	4.194.304	II22/	SHR	register(-ref.)-times a halvation
23	8.388.608	II22	OR	set bits set in of register(reference)
24	16.777.216	II22&	AND	clear bits not set in register(-ref.)
25	33.554.432	II21?	BTST	check if reg.(-ref.)-th bit is set
26	67.108.864	II21?	CMP	compare reg.(-ref.)1 with reg.(-ref.)2
27	134.217.728	:P00.	JMP	add integer to PC _u /EIP _u (=jump to)
28	268.435.456	CP1.<	JLT	jump if CMP <
29	536.870.912	CP1!>	JLE	jump if CMP ≤
30	1.073.741.824	CP1.=	JEQ	jump if CMP =
31	2.147.483.648	CP1!<	JGE	jump if CMP ≥
32	4.294.967.296	CP1!=	JNE	jump if CMP ≠
33	\$2.0000.0000	CP1.>	JGT	jump if CMP >
34	\$4.0000.0000	CP1!<	JPL	jump if ≥ 0
35	\$8.0000.0000	CP1.<	JMI	jump if < 0
36	\$10.0000.0000	CP1.^	JCS	jump if carry-flag is set
37	\$20.0000.0000	CP1!^	JCC	jump if carry-flag is clear
38	\$40.0000.0000	CP1.-	JVS	jump if overflow is set
39	\$80.0000.0000	CP1!~	JVC	jump if overflow is clear
40	\$100.0000.0000	CP2.<	DJMP	decrement and jump if reg.(-ref.) < 0
41	\$200.0000.0000	PS1..	CALL	PC _u /EIP _u →-(USP _u /ESP _u); + JUMP
42	\$400.0000.0000	SP11.	RET	(USP _u /ESP _u) + →PC _u /EIP _u
43	\$800.0000.0000	.I...	I???	unknown integer-operation
44	\$1000.0000.0000	.F...	F???	unknown floating-point-operation

45	\$2000.0000.0000	FF09	FINIT	initialise floating-point-unit
46	\$4000.0000.0000	FI12	FIST	store float.point-reg. → integer-reg.
47	\$8000.0000.0000	IF12	FILD	load integer-reg. → floating-point-reg.
48	\$1.0000*2 ³²	IF22+	FIADD	floating-point add constant integer
49	\$2.0000*2 ³²	IF22-	FISUB	floating-point subtract const.integer
50	\$4.0000*2 ³²	IF22*	FIMUL	floating-point multiply const.integer
51	\$8.0000*2 ³²	IF22/	FIDIV	floating-point divide by const.integer
52	\$10.0000*2 ³²	IF21?	FCMP	float.pt.compare with integer → flags
53	\$20.0000*2 ³²	:F02	FLD#	load constant to floating-point-reg.
54	\$40.0000*2 ³²	.F12!	FABS	build floating-point amount
55	\$80.0000*2 ³²	FF12	FLD	copy floating-point register
56	\$100.0000*2 ³²	FF22+	FADD	add 2 floating-point-register
57	\$200.0000*2 ³²	FF22-	FSUB	subtract 2 floating-point-register
58	\$400.0000*2 ³²	FF22*	FMUL	multiply 2 floating-point-register
59	\$800.0000*2 ³²	FF22/	FDIV	divide one float.pt.reg. by another
60	\$1000.0000*2 ³²	.F12@	FSQRT	square-root of a floating-point-reg.
61	\$2000.0000*2 ³²	.F12@	FSIN	floating-point-sine
62	\$4000.0000*2 ³²	.F12@	FCOS	floating-point-cosine
63	\$8000.0000*2 ³²	.F12@	FATAN	floating-point arc-tangent
64	\$1*2 ⁴⁸	FF22*	FEXP2	y: = y*2 ^x (anykind of exp.-func.)
65	\$2*2 ⁴⁸	FF22/	FLOG2	y: = x*log ₂ y (any kind of logarithm)
66	\$4*2 ⁴⁸	FF21?	FCMP	compare 2 float.-point-reg. → Flags
67	\$8*2 ⁴⁸	\$I11	SMOV	move from a special Register
68	\$10*2 ⁴⁸	\$I11	MOVS	move to a special Register
...

Fig.4b

... using the following operation-type character-codes:

1st char. = source , 2nd char. = dest.:	? = unknown, ambitious letter for all possible following
	. = nothing
	: = constant number
	I = integer-register-[reference]-value
	F = floating-point register
	C = condition-code register (last byte in EFlags#/SR ₁₀)
	P = instruction-pointer / program-counter (EIP#/PC ₁₆)
	S = stack-pointer reference-value
	~ = comparison-operation → flags
	\$ = a special-register like flags-, control-, debug-, ... -reg.
	! = logical NOT for comparison in the 4th field
3rd char. = number of source-registers:	including destination reg., if it's used as source-reg. too
4th char. = number of dest.-registers:	including flags-register but without the instruction-pointer
5th char. = effect of arithmetic:	? = unknown
	. = none
	! = amount negation bitwise_inversion
	+ = addition
	- = subtraction
	* = multiplication
	/ = division
	% = rest of division
	= setting of bits (mostly bitwise OR)
	& = clearing of bits (mostly bitwise AND)
	@ = trigonometric- or exponential-function
	> = greater? (CC-flags dependant action)
	< = less? (CC-flags dependant action)
	= = equal? (CC-flags dependant action)
	^ = carry? (CC-flags dependant action)
	~ = overflow? (CC-flags dependant action)

Fig.4c

00704802 10000

OpCode-Register-Table: [ORT - by opcode icl, initial-conditions concerned registers and the effect]
column: datatype value-range meaning:

OpCode	(PK)	integer	0..2 ³² -1	complete instruction, truncated if > 4 bytes
IniConNr	(PK)	signed byte	-31..30	current number of the used initial conditions
Register_ID_dest	(PK)	signed byte	0..127	one by execution concerned destination-reg. (see RIT)
Register_ID_source	(PK)	signed byte	-10..127	-1 or one possible source-register (see RIT).
value before change		integer	0..2 ³² -1	register(reference)-value before it was changed
value after change		integer	0..2 ³² -1	register(reference)-value after changing
gradient if unsigned		signed byte	-128..127	before/after-gradient, when defined as unsigned
gradient if signed		signed byte	-128..127	before/after-gradient, when defined as signed
value_source		integer	0..2 ³² -1	value of a possible source-register(reference)
Operations_BitCode		number	0..2 ¹²⁸ -1	bitmask, which flags all possible operations between this Register_ID_dest / Register_ID_source-combination (p.e. 2+2=2*2 using same reg.'s). Values see CIT, calculation see fig.19.

Fig. 5

For every register- or register-reference-modification of the same opcode-execution one entry is generated, which gives information about the register-reference-values before and after the execution and further information about the degree of changing and with it a hint to a possible operation and a possible source-register which were used. (Packed-, nibble-, or BCD-operations are not considered.)

The last address-register is the stack-pointer. The last data-register is the "energy"-register.

An address-register can be every register which value can be a pointer to memory which destination can be accessed using this register as a reference.

Several registers can be modified simultaneous - therefore this additional 1:n table, where *Register_ID_dest* means the identity of the changed register. Sometimes many register-(references) could be the source for one operation - this quantum increases through the sum over all possible operations.

Therefore the following tables identify the used opcode and the concerned register(s):

OpCode-Learn-Table: [OLT - ascertained effect of the opcode using the concerning initial conditions]

column:	datatype	value-range	meaning:

OpCode	(PK)	integer	0..2 ³² -1	complete instruction, truncated if > 4 bytes
IniConNr	(PK)	signed byte	-31..30	number of used initial condition
active ChkSum corrupt		boolean	1 0	flag: checksum of active AC-program changed
inactive ChkSum corrupt		boolean	1 0	flag: checksum of inactive AC-program changed
Exception Vect changed		signed byte	-128..0	Register ID of the (first) overwritten exception-vector
multiple Exc Vect chg		boolean	1 0	more than one exception-vector was overwritten
Processor Mode Changed		boolean	1 0	flag: processor-mode changed (p.e. trace cleared)
Number of Exception		byte	0..N-1	exception-number [0:=no exception] (if 0=exc.:+1)
OpCode_length_or_jump		signed byte	-128..127	EIP _{PC_{II}} after execution - EIP _{PC_{II}} before execution - 128=\$FF=long back-jump, 127=\$7F=long forward j.
CCR before execution		byte	0..255	CC-flags, which could cause a jump.
Register changed BitCode		number	0..2 ²⁸ -1	$\mathcal{E} \ 2^{\text{ORT}} \cdot \text{Register_ID_dest} \ \forall \ \text{ORT}(\text{opcode}, \text{IniConNr})$
Register source BitCode		number	0..2 ²⁸ -1	$\mathcal{E} \ 2^{\text{ORT}} \cdot \text{Register_ID_source} \ \forall \ \text{ORT}(\text{opcode}, \text{IniConNr})$
max_Operations_BitCode		number(19)	0..2 ²⁸ -1	$\mathcal{E} \ \text{ORT} \cdot \text{Calculation_BitCode} \ \forall \ \text{ORT}(\text{opcode}, \text{IniConNr})$
min_Operations_BitCode		number(19)	0..2 ²⁸ -1	$\mathcal{E} \ \text{ORT} \cdot \text{Calculation_BitCode} \ \forall \ \text{ORT}(\text{opcode}, \text{IniConNr})$
time of execution		integer	0..2 ³² -1	deci-seconds after {20.9.1994, 0:00:00, 0 o'clock}
cycles of execution		byte	1..255	clock-cycles the opcode-execution needed
aim_valuation		signed byte	-128..127	aim-attaining-valuation value above initial-conditions
gradient aim valuation		signed byte	-128..127	- difference to $CLT \ n-1, \text{IniConNr}, \text{aim_valuation}$

Fig. 6

OpCode-Base-Table: [OBT - ascertains the effect of the opcode-execution from the initial conditions]
 column: datatype value-range meaning:

OpCode (PK)	integer	0..2 ³² -1	complete instruction, truncated if > 4 bytes
Execution counter	byte	0..255	number of the OLT-entries until now
FatalError_counter	byte	0..255	number of the opcode-caused fatal errors until now: Checksum_corrupt, Exception_Vect_changed, Trace_Bit_cheared, Processor_Mode_changed, and the exceptions without Divide_Error, Overflow.
low_Error counter	byte	0..255	number of divide-errors plus overflow-exceptions
Jump longOp probability	signed byte	-128..127	probability that it's a long opcode or a jump
avg OpCpde jump length	signed byte	-128..127	average length of opcode or jump
OpCode len unconfirmed	boolean	1 0	min. one divergence from above average exists
avg cycles of execution	byte	1..255	average by opcode-execution needed clock-cycles
exec_cycles unconfirmed	boolean	1..0	min. one divergence from above average exists
Register write probability	signed byte	-128..127	probability: opcode writes into a register
Register copy probability	signed byte	-128..127	probability: opcode copies register
Memory write probability	signed byte	-128..127	probability: opcode writes into memory
Memory copy probability	signed byte	-128..127	probability: opcode copies memory
Reg to Mem probability	signed byte	-128..127	probability: opcode copies reg. to adr.reg.-destination
Mem to Reg probability	signed byte	-128..127	probability: opcode copies adr.reg.-destination to reg.
Multi Reg write prob	signed byte	-128..127	probability: opcode writes into more than one register
Multi Mem write prob	signed byte	-128..127	probability: opcode writes to many adr.reg.-destination
Multi Reg to Mem prob	signed byte	-128..127	prob.: opcode copies ≥2 reg. to ≥2 adr.reg.-destination
Multi Mem to Reg prob	signed byte	-128..127	prob.: opcode copies ≥2 adr.reg.-destinations to ≥2 reg
all_Reg_dest BitCode	number	0..2 ¹²⁸ -1	⊆ OLT.Register_changed BitCode ∨ OLT(OpCode)
cut_Reg_dest BitCode	number	0..2 ¹²⁸ -1	⊆ OLT.Register_changed BitCode ∨ OLT(OpCode)
all_Reg_source BitCode	number	0..2 ¹²⁸ -1	⊆ OLT.Register_source BitCode ∨ OLT(OpCode)
cut_Reg_source BitCode	number	0..2 ¹²⁸ -1	⊆ OLT.Register_source BitCode ∨ OLT(OpCode)
max_Operation BitCode	number	0..2 ¹²⁸ -1	⊆ OLT.max_Operation BitCode ∨ OLT(OpCode)
min_Operation BitCode	number	0..2 ¹²⁸ -1	⊆ OLT.min_Operation BitCode ∨ OLT(OpCode)
all_Operation BitCode	number	0..2 ¹²⁸ -1	⊆ OLT.min_Operation BitCode ∨ OLT(OpCode)
cut_Operation BitCode	number	0..2 ¹²⁸ -1	⊆ OLT.max_Operation BitCode ∨ OLT(OpCode)
max_write value	integer	0..2 ³² -1	maximum of all destination-values
min_write value	integer	0..2 ³² -1	minimum of all destination-values
avg_write value	integer	0..2 ³² -1	average over all destination-values
max_write gradient	integer	0..2 ³² -1	maximum gradient of the changed value
min_write gradient	integer	0..2 ³² -1	minimum gradient of the changed value
avg_write gradient	integer	0..2 ³² -1	average gradient of the changes value
evaluated_source Register	signed byte	-1, 0..127	ascertained source-register-ID (after OBT-evaluation)
evaluated_source_NumReg	signed byte	-128, -1, 0..127	-128 ≙ LOB means source-constant; 0..127 means a further source-register ID; none = -1 (after OBT-eval.)
evaluated_dest Register	signed byte	-1, 0..127	ascertained destination-register after OBT-evaluation
evaluated_dest_Register2	signed byte	-1, 0..127	possible 2nd dest.-reg. after OBT-evaluation or flags-reg. (2 real dest.-reg. ⇒ flags-reg. not appreciated).
evaluated_Operation_ID	signed byte	-1, 0..63	ascertained operation-ID (after OBT-evaluation)
Confirmation counter	byte	0..255	counter: same effects on other initial conditions
max_aim valuation	signed byte	-128..127	max. valuability of the opcode for aim-attaining
avg_aim valuation	signed byte	-128..127	average valuability of the opcode for aim-attaining
max_grad_aim valuation	signed byte	-128..127	max. gradient of aim-attaining in relation to the by 1 shorter opcode-combination of the last CBT(i-1).
avg_grad_aim valuation	signed byte	-128..127	average gradient concerning above -"

Fig. 7

Datatypes: Boolean 1 Bit, BCD/Nibble 4 Bit, Byte/char(1) 8 Bit, Word/short 16 Bit, DWord/Integer 32 Bit, QWord/number(19) 64 Bit, number/number(38,0) 128 Bit (38 digits ≙ 16 bytes), varchar2(N) string of variable length with max.N characters, long very long string with max(longDef) characters.

The following **combination-tables** are created dynamically - they have the same non-PK-columns, like the OBT respectively OLT respectively ORT, but for every additional number of combinations a further more opcode in the PK:

Combination-Register-Table: [CRT(i), i = number of opcodes in the combination, CRT(1) = ORT]
column: datatype value-range meaning:

OpCode 1 (PK)	integer	0-2 ³² -1	opcode #1 (first of the combination)
{for all OpCodes} (PK)	all integer	0-2 ³² -1	{for all opcodes from #2 up to N-1}
OpCode N (PK)	integer	0-2 ³² -1	opcode#N (last of the combination)
IniConNr (PK)	signed byte	-31..30	current number of the used initial conditions
Register ID dest (PK)	signed byte	0..127	one by execution concerned destination-reg. (see RIT)
Register ID source (PK)	signed byte	-1..127	-1 or one possible source-register (see RIT).
{same non-PK columns like in the Opcode-Register-Table.}	see above	see above	same non-PK columns like ORT.

Fig. 8

Combinations-Learn-Table: [CLT(i), i = number of opcodes in the combination, CLT(1) = OLT]
column: datatype value-range meaning:

OpCode 1 (PK)	integer	0-2 ³² -1	opcode #1 (first of the combination)
{for all OpCodes} (PK)	all integer	0-2 ³² -1	{for all opcodes from #2 up to N-1}
OpCode N (PK)	integer	0-2 ³² -1	opcode#N (last of the combination)
IniConNr (PK)	signed byte	-31..30	current number of the used initial conditions
{same non-PK columns like in the Opcode-Learn-Table.}	see above	see above	same non-PK columns like OLT.

Fig. 9

Combinations-Base-Table: [CBT(i), i = number of opcodes in the combination, CBT(1) = OBT]
column: datatype value-range meaning:

OpCode 1 (PK)	integer	0-2 ³² -1	opcode #1 (first of the combination)
{for all OpCodes} (PK)	all integer	0-2 ³² -1	{for all opcodes from #2 up to N-1}
OpCode N (PK)	integer	0-2 ³² -1	opcode#N (last of the combination)
{same non-PK columns like in the Opcode-Base-Table.}	see above	see above	same non-PK columns like OBT.

Fig. 10

CBT(max.) = CPT = Combination-Plan-Table = point of origin of the outcoming program.

Programming-aim and valuation-function tables:***Aim-Solution-Table: [AST - solutions of all programming-aims]***

column: datatype value-range meaning:

Aim ID (PK)	short	0..65535	Identifier of the programming-aim
Solution Nr (PK)	byte	0..255	number of the solution-program
aim_Program	long	String	opcode-combination of the solution here as a string
Program length	short	1..65535	length of the solution-program in doublewords
cycles of execution	integer	1..2 ³²⁻¹	execution-time in clock-cycles of the solution-program
used_Registers BitCode	number	1..2 ¹²⁸⁻¹	bitcode of all in the solution-program used registers
used_Operations BitCode	number	1..2 ¹²⁸⁻¹	bitcode of all in the solution-program used opcodes
used_aim_Valuation Func	signed short	0..32767	identifier of the used aim-distance valuation-function

Fig. 11

Aim-Description-Table: [ADT - identification and description of programming-aim]

column: datatype value-range meaning:

Aim ID (PK)	short	0..65535	Identifier of the programming-aim
aim_Description	varchar2(32)	≤32 Bytes	description of the programming-aim
used_Processor Mode	integer	0-2 ³²⁻¹	flags above CC control-register-bits
all_dest Register BitCode	number	1..2 ¹²⁸⁻¹	bitcode of all output-registers in this task
all_source Register BitCode	number	1..2 ¹²⁸⁻¹	bitcode of all input-registers in this task
unused_Register BitCode	number	1..2 ¹²⁸⁻¹	bitcode of all registers which should not be used
unused_Operation BitCode	number	0..2 ¹²⁸⁻¹	bitcode of all opcode-IDs which are not allowed to use in this task (default = \$0000.0000:0000.0000)
aim_implement_solutions	long	String	string of the Aim_ID's (words) of earlier solutions, which could be implemented here.
aim_fulfill_valuation mode	boolean	0 1	mode of aim-valuation: 0 = SQL ; 1 = machine-code
aim_fulfilled Flag Function	varchar2(99)	≤99 Bytes	boolean aim-attained recognition-function as a string
aim_Valuation FunctionID	signed short	0..32767	identifier of the valuation-function (see VFT)

Fig. 12

Functions-Identification-Table: [FIT - table of the basic subfunctions used in the valuation-function]

a.) for SQL-functions:

column: datatype value-range meaning:

Function ID (PK)	signed byte	-1..127	identification-number of the basic function
Function BitCode	number(19)	0..2 ⁶⁴⁻¹	bitcode of this basic function (only one bit is set)
Function Name	char(5)	5 Bytes	function-name
Function Type	byte	0..99	0 = value, 1 = unitary, 2 = binary, 3 = ternary, ...
Function Flatten	signed byte	-127..127	degree of flattening [+ = steepening, - = flattening]
Function_Template	varchar2(99)	≤99 Bytes	SQL function-template
Function_Description	varchar2(99)	≤99 Bytes	optional description of the basic sub-function

Fig. 13a

F.ID	Function BitCode	F.Name	F.T.	F.F.	Function Template	Function Description
0	1	NUM	0	0	< following value >	a constant number follows
1	2	ENGY	0	0	ELT.energy_after	energy after execution
2	4	GRAD	0	0	ELT.energy_after -ELT.energy_before	energy-gradient
3	8	VALUE	0	0	CLT(n). < columnNr >	value in the following column-number (last row)
4	16	EREG	0	0	< EnergyRegister_ID >	ID of the energy-register
5	32	SGN	1	0	SIGN(%s)	algebraic sign
6	64	ROUND	1	0	ROUND(%s, 0)	rounded
7	128	INT	1	0	FLOOR(%s)	truncated after dec.point
8	256	ABS	1	0	ABS(%s)	amount
9	512	NEG	1	0	-(%s)	negation
10	1.024	ADD	2	1	((%s) + (%s))	addition
11	2.048	SUB	2	-1	((%s) - (%s))	subtraction
12	4.096	MUL	2	4	((%s) * (%s))	multiplication
13	8.192	DIV	2	-4	((%s) / (%s))	division
14	16.384	MOD	2	-2	MOD(%s, %s)	rest of division
15	32.767	SQRT	1	-8	SQRT(%s)	square-root
16	\$1.0000	CBRT	1	-12	POWER(%s, 1/3)	cube-root
17	\$2.0000	MIN	2	-10	LEAST(%s, %s)	minimum
18	\$4.0000	MAX	2	-10	GREATEST(%s, %s)	maximum
19	\$8.0000	LN	1	-48	LN(%s)	natural logarithm
20	\$10.0000	EXP	1	48	EXP(%s)	nat. exponential-function
21	\$20.0000	LD	1	-32	LOG(2, %s)	logarithm on base 2
22	\$40.0000	POT2	1	32	POWER(2, %s)	2nd power of ...
23	\$80.0000	SIN	1	-64	SIN(%s)	sine
24	\$100.0000	COS	1	-64	COS(%s)	cosine
25	\$200.0000	TAN	1	127	TAN(%s)	tangent
26	\$400.0000	ASIN	1	127	ASIN(%s)	arc sine
27	\$800.0000	ACOS	1	127	ACOS(%s)	arc cosine
28	\$1000.0000	ATAN	1	-127	ATAN(%s)	arc tangent
29	\$2000.0000	SINH	1	40	SINH(%s)	sine hyperbolic
30	\$4000.0000	COSH	1	50	COSH(%s)	cosine hyperbolic
31	\$8000.0000	TANH	1	-127	TANH(%s)	tangent hyperbolic
32	\$1.0000.0000	LOG	2	-64	LOG(%s, %s)	logarithm
33	\$2.0000.0000	POT	2	64	POWER(%s, %s)	n-th power of ...
34	\$4.0000.0000	OR	2	1	((%s) (%s))	bitwise OR
35	\$8.0000.0000	AND	2	-1	((%s) & (%s))	bitwise AND
36	\$10.0000.0000	EQ	2	-127	DECODE(%s, %s, 1, 0)	equal
37	\$20.0000.0000	LE	2	-127	DECODE(GREATEST(%s - %s, 0), 0, 1, 0)	less-equal
38	\$40.0000.0000	GE	2	-127	DECODE(LEAST(%s - %s, 0), 0, 1, 0)	greater-equal
39	\$80.0000.0000	FRAME	1	-10	GREATEST(LEAST(%s, +127), -128)	frame to signed-byte: max. = 127, min. = -128
40	\$100.0000.0000	BITS	1	-64	(1 & %s) + (2 & %s) / 2 + (4 & %s) / 4 + (8 & %s) / 8 +	number of bits in the integer value
41	\$200.0000.0000	S_REG	0	0	ADT.all_source_Registers- BitCode	bitcode of the source- register
42	\$400.0000.0000	D_REG	0	0	ADT.all_dest_Registers BitCode	bitcode of dest.-register
43	\$800.0000.0000	AIM_F	0	0	VAL(ADT.aim_fulfilled_Flag- Function)	result of the boolean aim-fulfilled-function
...

Fig. 13b

b.) for machine-code functions:

column: datatype value-range meaning:

Function ID (PK)	signed byte	-1..127	identification-number of the basic function
Function BitCode	number(19)	0..2 ⁶⁴ -1	bitcode of this sub-function
Operations BitCode	number	0..2 ¹²⁸ -1	bitcode of the used opcodes in this function
Registers BitCode	number	0..2 ¹²⁸ -1	bitcode of the used registers in this function
Function Name	char(5)	5 Bytes	short notation of this sub-function
Function Type	byte	0..99	0 = value, 1 = unitary, 2 = binary, 3 = ternary, ...
Function Flatten	signed byte	-128..127	degree of function-flattening (1 \triangleq f(x) = x)
Function OpCodes	number	1..2 ¹²⁸ -1	sub-function in machine-code
Function Description	varchar2(99)	≤99 Bytes	optional description of the sub-function

Fig. 14a

Func.ID	Func.BitCode	Oper.BitCode	Reg.BitCode	Func.Name	F.T.	Func.OpCodes	Function Descript.
0	1	\$A000.4008	<energy>	FRAME	1	s.b. Func.1	prevent overflow
1	2	\$28800.0009	<energy>	SGN	1	s.b. Func.2	signum
2	4	\$0000.0002	<energy>	NEG	1	<NEG>	negation
3	8	\$0000.0200	<energy>	MUL2	1	<SHLI>	division by 2
4	16	\$0000.0400	<energy>	DIV2	1	<SRI>	multiplication by 2
5	32	\$0000.0100: 4A00.8018	<D0> (en)	ILOG2	1	s.b. Func.3	mogarithm dualis
6	64	\$1000.C000: 0000.0000	<FP0> (en)	ISQRT	1	s.b. Func.4	square-root
7	128		s. 1.4.2	ICBRT	1	s.above 1.4.2	cube-root
8	256	\$0000.8000	<en-1> (en)	MOV	2	<MOV>	copying of one reg. before energy-reg.
9	512	\$0000.8000	<en-1> (en)	SWAP	2	s.b. Func.5	swap with reg. before energy-reg.
10	1024	\$0001.0000	<en-1> (en)	ADD	2	<ADD>	addition with "-"
11	2048	\$0002.0000	<en-1> (en)	SUB	2	<SUB>	subtraction of "-"
12	4096	\$0004.0000	<en-1> (en)	MUL	2	<MUL>	multiplied with "-"
13	8192	\$0008.0000	<en-1> (en)	DIV	2	<DIV>	division by "-"
...

Fig. 14b

Function:	OpCodes of: (machine-code compilation of these mnemonics, here a Motorola-Example)
Func.1	CMPI 127,(E); JLE (+2); MOVI #127,(E); CMPI -128,(E); JGE (+2); MOVI #-128,(E)
Func.2	TST (E); JGE (+3); MOVI #-1,(E); JMP (+5); JGT (+3); MOVI #0,(E); JMP (+2); MOVI #+1,(E)
Func.3	MOVI #31,D0; BTST D0,(E); JEQ (+3); DJMP D0,(-2); ADDI #1,D0; MOVE D0,(E)
Func.4	FILD (E); FSRQT; FIST (E)
Func.5	MOVE (E-1),-(A7); MOVE (E),(E-1); MOVE (A7)+,(E)

Fig. 14c

00704802 110300 000011 008402 60

Valuation-Function-Table: [VFT - Table of the valuation-functions]

column: datatype value-range meaning:

Valuation_Function_ID (PK)	signed short	± 32767	identifier of the valuation-function (energyspecif. neg.)
Valuation_Function_Type	char(1)	'E' 'A'	'E' = energy-valuation, 'A' = valuation for reaching closeness to programming-aim, ... (maybe further)
Valuation_Function_Mode	boolean	0 1	0=SQL-mode; 1 = machine-code mode
Valuation_Function	varchar2(99)	≤99 Bytes	valuation-function for energy- or aim-attainment
execution_counter	integer	0-2 ³² -1	number of uses of this valuation-function
used_Functions_BitCode	number(19)	0..2 ⁶⁴ -1	BitCodes of all subfunctions
Function_ID_Chain	varchar2(99)	≤99 Bytes	chain of subfunctions (one byte ≙ one Function_ID)
avg_Func_execution_time	integer	0-2 ³² -1	average by val.-func.-execution needed clock-cycles
boundary_value_counter	integer	0-2 ³² -1	counter incremented if the result is -128 or +127
low_value_counter	integer	0-2 ³² -1	counter incremented if the result inside ±16
Valuation_Function_value	signed byte	-128..127	valuability of the valuation-function = SAC.Self-Valuation Aim/Energy(Valuation_Function, Values)

Fig. 15a Initial Entries for Energy-Valuation and Aim-Closeness-Valuation:

ID	Ty	M	Valuation_Function
-1	'E'	0	MAX(MIN SGN(EnergyReg' - EnergyReg °) · SQRT(EnergyReg' - EnergyReg °) - 32 · \sum CLT(i).Register_changed_BitCode & (1 2' Energy Register_ID) , +127), -128)
0	'A'	0	MAX(MIN(16 · \sum CLT(i).Register_changed_BitCode & ADT.all_dest_Register_BitCode } + 16 · \sum CLT(i).Register_source_BitCode & ADT.all_source_Register_BitCode } + 32 · ADT.aim_fulfilled_Flag_Function(Aim_ID) - CLT(i).Processor_Mode changed - ¼ · CLT(i).cycles_of_execution - (CLT(i).active inactive_ChkSum_corrupt) - (CLT(i).Exception_vect_changed > 0) - (CLT(i).Number_of_Exception > 0) - ½ (CLT(i).OpCode_length_or_jump > 4 or ≤ 0 , +127), -128)

ex#	used F. BitCode	Function_ID chain	ex.T	bdy#	low#	F.Val
0	\$189.0040.983B	2,5; 2,15; 12; 4,22,3,11,35,40,1,32,12; 11; 39	0	0	0	0
0	\$EE9.0001.3AAA	3,11,42,35,1,16,12; 3,12,41,35,1,16,12,10; 43,1,32,10, 3,7,11; 3,16,1,5,13,11; 3,3,11; 3,5,11 3,5,5,10; 3,8,5,10; 3,9,1,0,37,11;3,9,1,5,38,11;39	0	0	0	0

Fig. 15b**Status of the Artificial Consciousness:** [SAC - status-values of the AC-program (only 1 row)]

column: datatype value-range meaning:

Programm_StartDate	timestamp	datetime	date and time of the start of the AC-program
actual_Processor_Mode	integer	0-2 ³² -1	flags above CCR control-register-bits
actual_CPT_index	byte	1..255	CBT(max(i)=actual_CPT_Nr) = actual CPT
CxT_counter	short	1..65535	number of creations of the dynamic CxT-tables
Aims_total	short	1..65535	number of total programming-aims
Aims_solved	short	0..65535	number of solved programming-aims
actual_Aim_ID	short	0..65535	ID of the actual programming-aim
Aim_Valuation_Mode	boolean	0 1	mode of the programming-aim-specific valuation-function: 0=SQL-mode; 1 = machine-code mode
Aim_Valuation_FunctionID	signed short	0..32767	actual VFT.Valuation_Function_ID referring the closeness to the programming-aim
Aim_Self_Valuation_Func	varchar2 (400)	max. 400 Chars.	PL/SQL-valuation-function referring the efficiency of the valuation-function
Energy_Valuation_Mode	boolean	0 1	mode of the energyspecific valuation-function 0=SQL-mode; 1 = machine-code mode
Energy_Valuation_Func_ID	signed short	-1..-32768	actual VFT.Valuation_Function_ID for energy-valuation
Energy_Self_Valuation_Func	varchar2 (400)	max. 400 Chars.	PL/SQL-valuation-function referring the efficiency of the energyspecific valuation-function
max_Valuation_Function	signed short	0..32767	highest ID of all valuation-functions in the VFT.
min_Valuation_Function	signed short	-1..-32768	lowest ID of all valuation-functions in the VFT.

Fig. 16

Energy-Learn-Table: [ELT - appraises energyspecific actions in dependence of the used initial-conditions]
 column: datatype value-range meaning:

Energy_action (PK)	number	0..2 ¹²⁸⁻¹	max.16 byte opcode-combination of the action which changed the energy-register.
IniConNr (PK)	signed byte	-31..30	number of the used initial condition
Energy_before	integer	0..2 ³²⁻¹	energy-register before execution
Energy_after	integer	0..2 ³²⁻¹	energy-register after execution
min_Operations BitCode	number	0..2 ¹²⁸⁻¹	bitcode of the probably used opcodes.
max_Operations BitCode	number	0..2 ¹²⁸⁻¹	bitcode of the possibly used opcodes.
Register_changed BitCode	number	1..2 ¹²⁸⁻¹	bitcode of the by action changed registers.
Register_source BitCode	number	1..2 ¹²⁸⁻¹	bitcode of the probable source-registers.
used_cycles of execution	short	1..65535	needed clock cycles for the energyspecific action
Energy_valuation	signed byte	-128..127	result of the actual VFT.Energy_valuation_Function
Valuation_Function_ID	signed short	-1..-32768	used energyspecific valuation-function

Fig. 17

Energy-Base-Table: [EBT - evaluation of the energyspecific actions]
 column: datatype value-range meaning:

Energy_action (PK)	number	0..2 ¹²⁸⁻¹	max.16 byte opcode-combination of the action which changed the energy-register.
Execution_counter	byte	0..255	number of the ELT-entries until now.
FatalError_counter	byte	0..255	number of the occurred fatal errors: fatal errors correlate the columns 3-7 of the above learn-table, except Divide-Error or Overflow-Exc.
low_Error_counter	byte	0..255	number of Divide-Errors or Overflow - Exceptions
avg_Energy_after	integer	0..2 ³²⁻¹	average energy-value after the action
all_Reg_dest BitCode	number	0..2 ¹²⁸⁻¹	⌈ ELT.Register_changed BitCode ∨ ELT(OpCode)
cut_Reg_dest BitCode	number	0..2 ¹²⁸⁻¹	⌋ ELT.Register_changed BitCode ∨ ELT(OpCode)
all_Reg_source BitCode	number	0..2 ¹²⁸⁻¹	⌈ ELT.Register_source BitCode ∨ ELT(OpCode)
cut_Reg_source BitCode	number	0..2 ¹²⁸⁻¹	⌋ ELT.Register_source BitCode ∨ ELT(OpCode)
max_Operation BitCode	number	0..2 ¹²⁸⁻¹	⌈ ELT.max_Operation BitCode ∨ ELT(OpCode)
min_Operation BitCode	number	0..2 ¹²⁸⁻¹	⌋ ELT.min_Operation BitCode ∨ ELT(OpCode)
all_Operation BitCode	number	0..2 ¹²⁸⁻¹	⌈ ELT.min_Operation BitCode ∨ ELT(OpCode)
cut_Operation BitCode	number	0..2 ¹²⁸⁻¹	⌋ ELT.max_Operation BitCode ∨ ELT(OpCode)
max_write_value	integer	0..2 ³²⁻¹	maximum of all energy-values after energy-action
min_write_value	integer	0..2 ³²⁻¹	minimum of all energy-values after energy-action
avg_write_value	integer	0..2 ³²⁻¹	average of all energy-values after energy-action
max_write_gradient	integer	0..2 ³²⁻¹	maximum gradient of the changes energy-register
min_write_gradient	integer	0..2 ³²⁻¹	minimum gradient of the changes energy-register ##
avg_write_gradient	integer	0..2 ³²⁻¹	average gradient of the changes energy-register
equal_value_probability	signed byte	-128..127	probability of equal result of energyspecific action
avg_Energy_gradient	signed int	± 2 ³¹	average value-gradient of this energyspecific action
equal_Gradient_probability	signed byte	-128..127	probability: gradient is constant
avg_cycles of execution	short	1..65535	average needed clock cycles for this action
avg_Energy_valuation	signed byte	-128..127	result of the actual VFT.Energy_valuation_Function
Valuation_Function_ID	signed short	-1..-32768	ID of the used energy-valuation-function

Fig. 18

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3.2 Flowchart of the AC-Program:

3.2.1 CxT(i) value assignments:

ORT & CRT(i) value assignments:

ORT.Register_ID_dest := $\log_2(\text{Bit}(\text{OLT.Register_changed_Mask}), \text{of the regarded changing})$
ORT.Register_ID_source := Register_ID(C°), if ORT.calculation code > 0, otherwise -1
ORT.value before change := value(Register_ID_dest), before opcode-execution
ORT.value after change := value(Register_ID_dest), after opcode-execution
ORT.gradient if signed := $\text{MAX}[\text{MIN}(\text{ORT.value after change} - \text{ORT.value before change}, +127), -128]$
ORT.gradient if unsigned := $\text{MAX}[\text{MIN}(\text{ORT.value after change} - \text{ORT.value before change}, +127), 128]$
ORT.Operation_BitCode := $1 - (\text{Flags} \neq \text{Flags}^\circ) \&\&\vee (V = V^\circ) \&\&[\text{NF} \&\& (V_1^\circ < 0) \mid \mid \text{ZF} \&\& (V_1^\circ = 0)]$ $+ 2 \cdot [(V_1' = -V_1^\circ) \&\&\vee (V = V^\circ)] + 4 \cdot [(V_1' = -V_1^\circ) \&\&\vee (V = V^\circ)] + 8 \cdot [(V_1' = 0\text{LB}) \&\&\vee (V = V^\circ)]$ $+ 16 \cdot [(V_1' = V_1^\circ + 0\text{LB}) \&\&\vee (V = V^\circ)] + 32 \cdot [(V_1' = V_1^\circ - 0\text{LB}) \&\&\vee (V = V^\circ)] + 64 \cdot [(V_1' = V_1^\circ - 0\text{LB}) \&\&\vee (V = V^\circ)]$ $+ 128 \cdot [(V_1' = V_1^\circ / 0\text{LB}) \&\&\vee (V = V^\circ)] + 256 \cdot [(V_1' = V_1^\circ \% 0\text{LB}) \&\&\vee (V = V^\circ)] + 512 \cdot [(V_1' = V_1^\circ \cdot 2^\circ - 0\text{LB}) \&\&\vee (V = V^\circ)]$ $+ 2^{10} \cdot [(V_1' = V_1^\circ / 2^\circ - 0\text{LB}) \&\&\vee (V = V^\circ)] + 2^{11} \cdot [(V_1' = V_1^\circ \mid 0\text{LB}) \&\&\vee (V = V^\circ)] + 2^{12} \cdot [(V_1' = V_1^\circ \& 0\text{LB}) \&\&\vee (V = V^\circ)]$ $+ 2^{13} \cdot (\text{Flags} \neq \text{Flags}^\circ) \&\&\vee (V = V^\circ) \&\&[(\text{ZF} = 1) \&\& (2^\circ - 0\text{LB} \mid \sim V^\circ) \mid \mid (\text{ZF} = 0) \&\& (2^\circ - 0\text{LB} \mid V_1^\circ)]$ $+ 2^{14} \cdot (\text{Flags} \neq \text{Flags}^\circ) \&\&\vee (V = V^\circ) \&\&[\text{NF} \&\& (V_1^\circ < 0\text{LB}) \mid \mid \text{ZF} \&\& (V_1^\circ = 0\text{LB})] + 2^{15} \cdot [(V_1' = C_1^\circ) \&\&\vee (V = V^\circ)]$ $+ 2^{16} \cdot [(V_1' = V_1^\circ \cdot C_1^\circ) \&\&\vee (V = V^\circ)] + 2^{17} \cdot [(V_1' = V_1^\circ \cdot C_1^\circ) \&\&\vee (V = V^\circ)] + 2^{18} \cdot [(V_1' = V_1^\circ \cdot C_1^\circ) \&\&\vee (V = V^\circ)]$ $+ 2^{19} \cdot [(V_1' = V_1^\circ / C_1^\circ) \&\&\vee (V = V^\circ)] + 2^{20} \cdot [(V_1' = V_1^\circ \% C_1^\circ) \&\&\vee (V = V^\circ)] + 2^{21} \cdot [(V_1' = 2^{C_1^\circ} \cdot V_1^\circ) \&\&\vee (V = V^\circ)]$ $+ 2^{22} \cdot [(V_1' = V_1^\circ / 2^{C_1^\circ}) \&\&\vee (V = V^\circ)] + 2^{23} \cdot [(V_1' = V_1^\circ \mid C_1^\circ) \&\&\vee (V = V^\circ)] + 2^{24} \cdot [(V_1' = V_1^\circ \& C_1^\circ) \&\&\vee (V = V^\circ)]$ $+ 2^{25} \cdot (\text{Flags} \neq \text{Flags}^\circ) \&\&\vee (V = V^\circ) \&\&[(\text{ZF} = 1) \&\& (2^\circ \cdot C_1^\circ \mid \sim V_1^\circ) \mid \mid (\text{ZF} = 0) \&\& (2^\circ \cdot C_1^\circ \mid V_1^\circ)]$ $+ 2^{26} \cdot (\text{Flags} \neq \text{Flags}^\circ) \&\&\vee (V = V^\circ) \&\&[\text{NF} \&\& (V_1^\circ < C_1^\circ) \mid \mid \text{ZF} \&\& (V_1^\circ = C_1^\circ)]$ $+ 2^{27} \cdot [(IP' \leq IP^\circ) \mid \mid (IP' > IP^\circ + 4)] \&\& (\text{Flags} = \text{Flags}^\circ) \&\&\vee (V = V^\circ)$ $+ 2^{28} \cdot [(IP' \leq IP^\circ) \mid \mid (IP' > IP^\circ + 4)] \&\& (\text{Flags} = \text{Flags}^\circ) \&\&\vee (V = V^\circ) \&\&[\text{NF} \&\& \text{IF} \mid \mid \text{NF} \&\& \text{VF} \mid \mid \dots \vee \text{Jcc}(\text{CCR})]$ $+ 2^{40} \cdot [(IP' \leq IP^\circ) \mid \mid (IP' > IP^\circ + 4)] \&\& (V_1' = V_1^\circ - 1) \mid \mid (V_1' = -1) \&\& (\text{Flags} = \text{Flags}^\circ) \&\&\vee (V = V^\circ)$ $+ 2^{41} \cdot [(IP' = IP^\circ \pm 0\text{LB}) \&\& (SP) = IP^\circ] \&\& (\text{Flags} = \text{Flags}^\circ) \&\&\vee (V = V^\circ)$ $+ 2^{42} \cdot [(IP' = -4(\text{SP})) \&\& (\text{Flags} = \text{Flags}^\circ) \&\&\vee (V = V^\circ) + 2^{43} \cdot [(V_1' \neq V_1^\circ) \&\& (! \text{other_Integer_Operation_BitCode})]$ $+ 2^{44} \cdot [(V_1' \neq V_1^\circ) \&\& (! \text{other_FloatingPoint_Operation_BitCode})] + 2^{45} \cdot [(\text{CCR-Flags} = 0) \&\& (V_1' = 0)]$ $+ 2^{46} \cdot [(V_1' = C_F^\circ) \&\&\vee (V = V^\circ)] + 2^{47} \cdot [(V_1' = C_1^\circ) \&\&\vee (V = V^\circ)] + 2^{48} \cdot [(V_1' = V_F^\circ + C_1^\circ) \&\&\vee (V = V^\circ)]$ $+ 2^{49} \cdot [(V_1' = V_F^\circ - C_1^\circ) \&\&\vee (V = V^\circ)] + 2^{50} \cdot [(V_1' = V_F^\circ \cdot C_1^\circ) \&\&\vee (V = V^\circ)] + 2^{51} \cdot [(V_1' = V_F^\circ / C_1^\circ) \&\&\vee (V = V^\circ)]$ $+ 2^{52} \cdot (\text{Flags} = \text{Flags}^\circ) \&\&\vee (V = V^\circ) \&\&[\text{NF} \&\& (V_F^\circ < C_1^\circ) \mid \mid \text{ZF} \&\& (V_F^\circ = C_1^\circ)]$ $+ 2^{53} \cdot [(V_F' = 1.0) \mid \mid (V_F' = 0.0) \mid \mid (V_F' = \pi) \mid \mid (V_F' = e)] \&\&\vee (V = V^\circ)$ $+ 2^{54} \cdot [(V_F' = -V_F^\circ) \&\& (V_F^\circ < 0) \&\&\vee (V = V^\circ)] + 2^{55} \cdot [(V_F' = C_F^\circ) \&\&\vee (V = V^\circ)]$ $+ 2^{56} \cdot [(V_F' = V_F^\circ + C_F^\circ) \&\&\vee (V = V^\circ)] + 2^{57} \cdot [(V_F' = V_F^\circ - C_F^\circ) \&\&\vee (V = V^\circ)]$ $+ 2^{58} \cdot [(V_F' = V_F^\circ \cdot C_F^\circ) \&\&\vee (V = V^\circ)] + 2^{59} \cdot [(V_F' = V_F^\circ / V_F^\circ) \&\&\vee (V = V^\circ)] + 2^{60} \cdot [(V_F' \cdot V_F^\circ = V_F^\circ) \&\&\vee (V = V^\circ)]$ $+ 2^{61} \cdot [V_F' = \sin(V_F^\circ) \&\&\vee (V = V^\circ)] + 2^{62} \cdot [V_F' = \cos(V_F^\circ) \&\&\vee (V = V^\circ)] + 2^{63} \cdot [(V_F' = \text{atan}(V_F^\circ)) \&\&\vee (V = V^\circ)]$ $+ 2^{64} \cdot [(V_F' = V_F^\circ \cdot 2^\circ \cdot V_{F-1}^\circ) \&\&\vee (V = V^\circ)] + 2^{65} \cdot [(V_F' = V_{F-1}^\circ \cdot \log_2(V_F^\circ) \&\&\vee (V = V^\circ)]$ $+ 2^{66} \cdot (\text{Flags} = \text{Flags}^\circ) \&\&\vee (V = V^\circ) \&\&[\text{NF} \&\& (V_F^\circ < C_F^\circ) \mid \mid \text{ZF} \&\& (V_F^\circ = C_F^\circ)] + 2^{67} \cdot [(V_1' = C_8^\circ) \&\&\vee (V = V^\circ)]$ $+ 2^{68} \cdot [(V_8' = C^\circ) \&\&\vee (V = V^\circ)] + \dots$, where $V' = \text{value_after_change}(\neg \text{Flags})$, $V^\circ = \text{value_before_change}$ $C^\circ = \text{value}(\text{Register_ID_source})$. Here has to be checked over all Register_ID_source(eq.kind). Though equal Register_ID's in the PK several bits can be set. [p.e. because $4 = 2 + 2 = 2 \cdot 2 = \text{SHL}(2) = \dots$]

Fig. 19

OLT & CLT(i) value assignments:

OLT.Processor_Mode_Changed := $\neg \{ \text{EFlags}_{\text{SR}_{i-1}} \mid \neg 2^\circ \text{CCR_Flags} \} > 0 \mid \mid$ ORT.value after changel Register_ID of a special register)
OLT.aim_valuation := VFT.Aim_Valuation_Function(SAC.Aim_Valuation_FunctionID, ORT.xxxxx, Registers_changed_BitCode, Registers_source_BitCode, min_Operations_BitCode, max_Operations_BitCode, used_cycles_of_execution, ...)
CLT(n).gradient_aim_valuation := CLT(n).aim_valuation - CLT(n-1).aim_valuation
all other column-assignments are declared adequate in the OLT-description in fig.6.

Fig. 20

OBT & CBT(i) value-assignments:

OBT.Execution_counter := Execution_counter + 1
OBT.FatalError_counter := FatalError_counter + (0 < OLT.Number_of_Exception ≠ Divide_Error, Overflow) OLT.active_ChkSum_corrupt OLT.inactive_ChkSum_corrupt OLT.Exception_vect_changed OLT.Processor_Mode_changed)
OBT.Jump_longOp_probability := MAX(MIN(Jump_probability + (OLT.OpCode_length_or_jump ≤ 0) + (OLT.OpCode_length_or_jump > 4), +127], -128]
OBT.avg_OpCode_jump_length := (execution_counter * avg_OpCode_jump_length + akt.OpCode_jump_length) / (execution_counter + 1)
OBT.OpCode_len_unconfirmed := OpCode_len_unconfirmed (avg_OpCode_length ≠ act.OpCode_length)
OBT.avg_cycles_of_execution := (execution_counter * avg_cycles_of_execution + act.cycles_of_execution) / (execution_counter + 1)
OBT.exec_cycles_unconfirmed := exec_cycles_unconfirmed (avg_cycles_of_execution ≠ act.cycles_of_execution)
OBT.Register_write_probability := MAX(MIN(Register_write_probability + 2 * ((min.Reg.ID ≤ ORT.Column_ID_OLT ≤ max.Reg.ID) && ORT.value_before_change ≠ ORT.value_after_change] - 1, +127], -128]
OBT.Register_copy_probability := MAX(MIN(MIN(Register_copy_probability + 2 * ((min.Reg.ID ≤ ORT.Column_ID_OLT ≤ max.Reg.ID) && ORT.value_before_change ≠ ORT.value_after_change && (min.Reg.ID ≤ ORT.Column_ID_source ≤ max.Reg.ID)] - 1, +127], -128]
OBT.Memory_write_probability := MAX(MIN(Memory_write_probability + 2 * ((min.Adr.Reg.ID ≤ ORT.Column_ID_OLT ≤ max.Adr.Reg.ID) && ORT.value_before_change ≠ ORT.value_after_change] - 1, +127], -128]
OBT.Memory_copy_probability := MAX(MIN(Memory_copy_probability + 2 * ((min.Adr.Reg.ID ≤ ORT.Column_ID_OLT ≤ max.Adr.Reg.ID) && ORT.value_before_change ≠ ORT.value_after_change && (min.Adr.Reg.ID ≤ ORT.Column_ID_source ≤ max.Adr.Reg.ID)] - 1, +127], -128]
OBT.Reg_to_Mem_probability := MAX(MIN(Reg_to_Mem_probability + 2 * ((min.Adr.Reg.ID ≤ ORT.Column_ID_OLT ≤ max.Adr.Reg.ID) && ORT.value_before_change ≠ ORT.value_after_change && (min.Reg.ID ≤ ORT.Column_ID_source ≤ max.Reg.ID)] - 1, +127], -128]
OBT.Mem_to_Reg_probability := MAX(MIN(Mem_to_Reg_probability + 2 * ((min.Reg.ID ≤ ORT.Column_ID_OLT ≤ max.Reg.ID) && ORT.value_before_change ≠ ORT.value_after_change && (min.Adr.Reg.ID ≤ ORT.Column_ID_source ≤ max.Adr.Reg.ID)] - 1, +127], -128]
OBT.Multi_Reg_write_prob := <i>like in Register_write_probability, but with min.2 appropriate ORT.Column_ID_OLT-entries.</i>
OBT.Multi_Mem_write_prob := <i>like in Memory_write_probability, but with min.2 appropriate ORT.Column_ID_OLT-entries.</i>
OBT.Multi_Reg_to_Mem_prob := <i>like in Reg_to_Mem_probability, but with min.2 appropriate ORT.Column_ID_OLT + Column_ID_source-entries.</i>
OBT.Multi_Mem_to_Reg_prob := <i>like in Mem_to_Reg_probability, but with min.2 appropriate ORT.Column_ID_OLT + Column_ID_source-entries.</i>
OBT.xxx_Reg_source dest_BitCode: <i>see table-description</i>
OBT.xxx_calculation_BitCode: <i>see table-description</i>
OBT.max_write_value := MAX(max_write_value, ORT.value_after_change)
OBT.min_write_value := MIN(min_write_value, ORT.value_after_change)
OBT.avg_write_value := (execution_counter * avg_write_value + ORT.value_after_change) / (execution_counter + 1)
OBT.max_write_gradient := MAX(max_write_gradient, ORT.value_after_change - ORT.value_before_change)
OBT.min_write_gradient := MIN(min_write_gradient, ORT.value_after_change - ORT.value_before_change)
OBT.avg_write_gradient := (execution_counter * avg_write_gradient + ORT.value_after_change - ORT.value_before_change) / (execution_counter + 1)
OBT.evaluated_source [Num]Register := <i>probability-function</i> (xxx_Reg_source_BitCode, confirmation_counter)

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OBT.evaluated dest Register[2] := <i>probability-function</i> (xxx Reg_dest BitCode, confirmation-counter)
OBT.evaluated Operation ID := <i>probability-function</i> (xxx Operation BitCode, confirmation-counter)
OBT.Confirmation_counter := Confirmation_counter + <i>exist(equivalent OLT+ORT-entry with lower IniConNr)</i>
OBT.max_aim_valuation := MAX(max_aim_valuation, OLT.aim_valuation)
OBT.avg_aim_valuation := (execution_counter * avg_aim_valuation + OLT.aim_valuation) / (execution_counter + 1)
CBT(n).max_grad_aim_valuation := MAX(CBT(n).max_aim_valuation, CLT(n).aim_valuation) - CBT(n-1).max_aim_valuation.
CBT(n).avg_grad_aim_valuation := (execution_counter * CBT(n).avg_aim_valuation + CLT(n).aim_valuation) / (execution_counter + 1) - CBT(n-1).avg_grad_aim_valuation

Fig.21

3.2.2 ELT and EBT value-assignments:

ELT.max Operations BitCode := OLT.max Operations OpCode
ELT.min Operations BitCode := OLT.min Operations OpCode
ELT.Register changed BitCode := OLT.Registers changed BitCode
ELT.Register source BitCode := OLT.Registers source BitCode
ELT.Energy_Valuation := VFT.Energy_valuation_Function(SAC.Energy_Valuation_FunctionID, Energy_after, Energy_before, Registers_changed_BitCode, Registers_source_BitCode, min_Operations_BitCode, max_Operations_BitCode, used_cycles_of_execution, ...)
ELT.Valuation Function ID := <i>for calculation of Energy Valuation used VFT.Valuation Function ID</i>
EBT.avg_Energy_after := (execution_counter * avg_Energy_after + ELT.Energy_after) / (execution_counter + 1)
EBT.equal value probability := equal value probability + 2 * (avg_Energy_after = ELT.Energy_after) - 1
EBT.avg_Energy_gradient := (execution_counter * avg_Energy_gradient + ELT.Energy_after - ELT.Energy_before) / (execution_counter + 1)
EBT.equal gradient probability := equal gradient probability + 2 * (avg_Energy_gradient = ELT.Energy_after - ELT.Energy_before) - 1
EBT.xxx Operations/Registers BitCode <i>see table-description</i>
EBT.avg_cycles_of_execution := (execution_counter * avg_cycles_of_execution + ELT.used_cycles_of_execution) / (execution_counter + 1)
EBT.avg_Energy_Valuation := (execution_counter * avg_Energy_Valuation + ELT.Energy_valuation) / (execution_counter + 1)

Fig.22

3.2.3 Definitions needed to read the flowchart:

- directives denotes a directive or a short sequence of directives.
 < *condition fulfilled ?* > Yes: branches horizontally, No: continue below.
 < *continuing-label* > denotes a label to or from another part of the flowchart.
block of directives denotes a block of earlier defined directives.

In the flowchart because of the complexity not all things are described until the smallest detail, but the fundamental functionality is presented clear and comprehensible.
 Self-evident things like closing a database-cursor or cofilling not explicit mentioned but existing table-columns (which do not need a special algorithm) are not performed additionally, because the meaning of these columns is already declared in 3.1.2 an their assignment-formulas in 3.2.1 or in 3.2.2.

In the flowchart means "generate ORT-entry and actualise OBT" what is already shown in the value-assignments in fig. 19-21.

Fig.23

3.2.4 AC-flowchart:

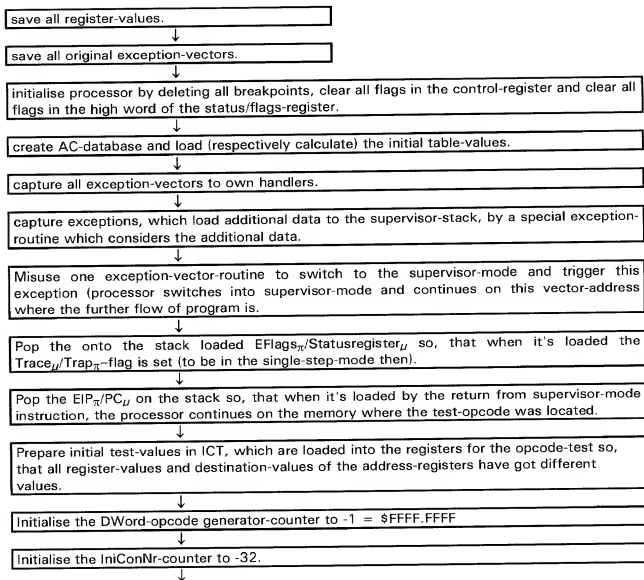
a.) Initial Preparations:

Fig.24a

00000000 00000000 00000000 00000000

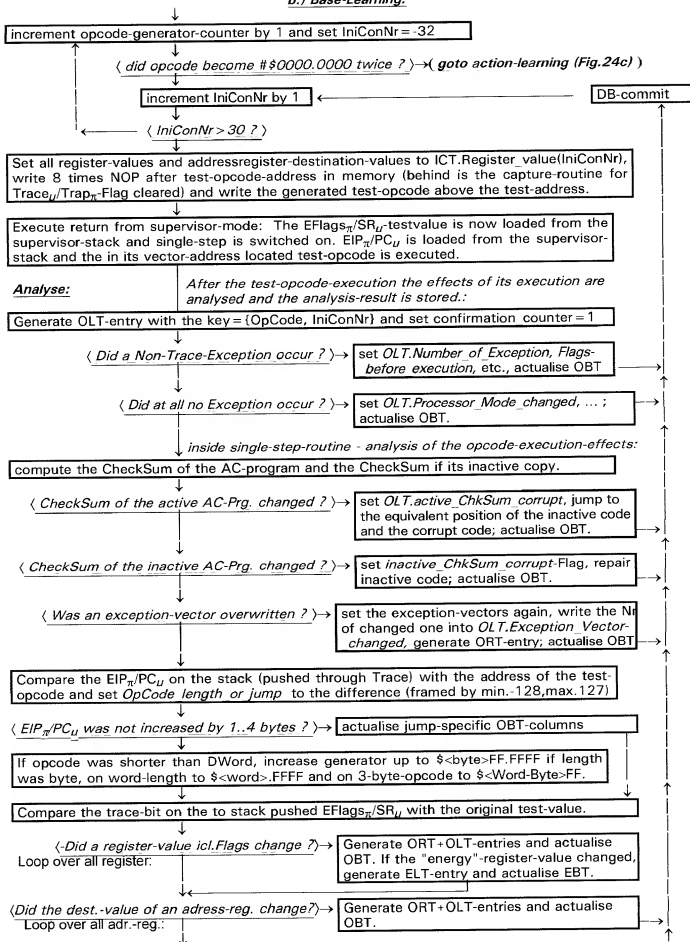
b.) Base-Learning:

Fig.24b

c.) Double-OpCode-Acting:**(Begin of Double-OpCode-Acting [after Base-Learning - Fig.24b])**

OBT-Evaluation: with probability-functions evaluated *source [Num]Register* is derived from OBT.xxx_Reg_source_BitCode , *evaluated dest Register[2]* derived from xxx_Reg_dest_BitCode and *evaluated Operation ID* is derived from xxx Operation BitCode (considering confirmation counter).

Define the last data-register as "energy"-register and set to an average value.

Open 1st Cursor over the OBT.

Fetch next row from the 1st OBT-cursor.

(Triple-Opcode-Planning (Fig.24c)) ← { Fetch1 empty (1st opcode processed) ? }

{ Jump_longOp_probability > 0 or Y[ADT.unused_Register_BitCode(Aim_ID) & (OBT.cut_source_Reg_BitCode|OBT.cut_dest_Reg_BitCode)] > 0 ? }

{ Execution_counter / OBT.OpCode_FatalError_counter(opcode1) < 5 ? }

Open 2nd Cursor over the OBT.

Fetch next row from the 2nd OBT-cursor.

{ Fetch2 empty (2nd opcode processed) ? }

{ Jump_longOp_probability > 0 or Y[ADT.unused_Register_BitCode(Aim_ID) & (OBT.cut_source_Reg_BitCode|OBT.cut_dest_Reg_BitCode)] (OpC.2) > 0 ? }

{ Execution_counter / OBT.OpCode_FatalError_counter(opcode2) < 5 ? }

initialise IniConNr = -32

IniConNr++

{ IniConNr > +30 ? }

Initialising and execution using the initial conditions of the corresponding IniConNr, like during the base-learning, but now for the double-opcode.

Same procedure like in the analysis-block of the base-learning (Fig.24b), but analysis-results now stored into CRT(2), CLT(2) and CBT(2).

decrement "energy"-register by 1.

{ Is the "energy"-register in the middle or high range ? }

Select and execute an EBT.Energy specific action which has a high energy action valuation and a high confirmation counter.

Analysis-block like above.

Fig.24c

d.) Triple-OpCode-Planning:

(Begin of Triple-OpCode-Planning (after Double-OpCode-Actions))

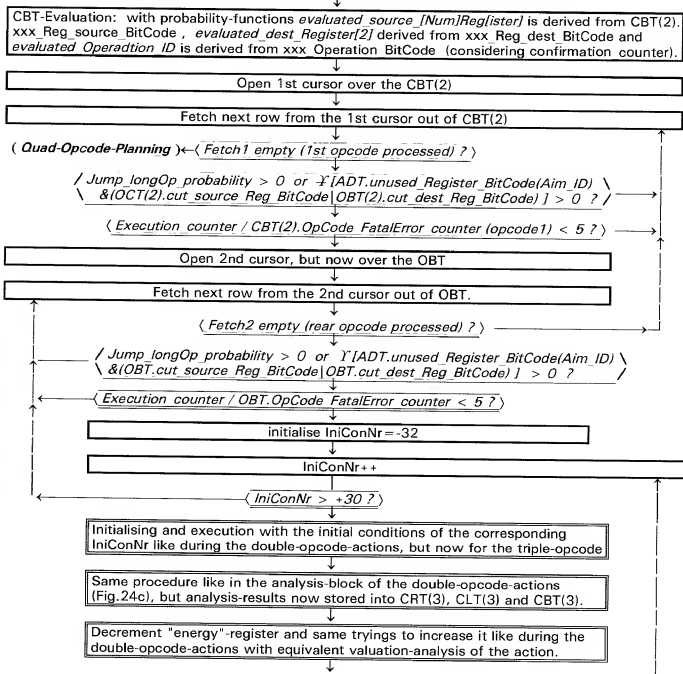


Fig.24d

Procedure for higher combinations analogous, using CxT(n), where n = sum_of_opcodes.

Good Morning

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Als nachstehend benannter Erfinder erkläre ich hiermit an Eides Statt:

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Method for generating a simple kind of Artificial Consciousness in a computer, which has the
capability to plan, generate automatically and execute machine-code for the solution of
arbitrary programming-abandonments. (Automatic Programming)

deren Beschreibung hier beigefügt ist, es sei denn (in diesem Falle Zutreffendes bitte ankreuzen), diese Erfindung

- ☐ wurde angemeldet am _____
unter der US-Anmeldenummer oder unter der
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Vertrags über die Zusammenarbeit auf dem Gebiet
des Patentwesens (PCT)
_____ und am
_____ abgeändert (falls
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Ich bestätige hiermit, daß ich den Inhalt der oben angegebenen Patentanmeldung, einschließlich der Ansprüche, die eventuell durch einen oben erwähnten Zusatzantrag abgeändert wurde, durchgesehen und verstanden habe.

Ich erkenne meine Pflicht zur Offenbarung jeglicher Informationen an, die zur Prüfung der Patentfähigkeit in Einklang mit Titel 37, Code of Federal Regulations, § 1.56 von Belang sind.

As a below named inventor, I hereby declare that:

My residence, post office address and citizenship are as stated next to my name.

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled

the specification of which is attached hereto unless the following box is checked:

- ☐ was filed on _____
as United States Application Number or PCT
International Application Number
_____ and was amended on
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I hereby state that I have reviewed and understand the contents of the above identified specification, including the claims, as amended by any amendment referred to above.

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Prior Foreign Applications
(Frühere ausländische Anmeldungen)

19352587.0-53 Germany
(Number) (Country)
(Nummer) (Land)

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(Nummer) (Land)

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(Application No.) (Filing Date)
(Ankennzeichen) (Anmeldetag)

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Priority Not Claimed
~~Priorität nicht beansprucht~~

02. Nov. 1999
(Day/Month/Year Filed)
(Tag/Monat/Jahr der Anmeldung)

(Day/Month/Year Filed)
(Tag/Monat/Jahr der Anmeldung)

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(Status) (patentiert, schwebend, aufgegeben)

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POWER OF ATTORNEY: As a named inventor, I hereby appoint the following attorney(s) and/or agent(s) to prosecute this application and transact all business in the Patent and Trademark Office connected therewith: (list name and registration number)

No Attorney

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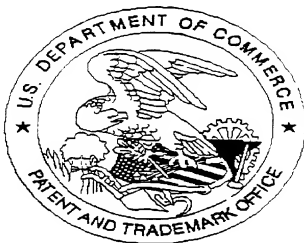
Direct Telephone Calls to: (name and telephone number)

Vor- und Zuname des einzigen oder ersten Erfinders <u>Gerd Krämer</u>	Full name of sole or first inventor <u>Gerd Krämer</u>
Unterschrift des Erfinders <u>G. Krämer</u> Datum <u>23.10.00</u>	Inventor's signature <u>G. Krämer</u> Date <u>23.10.00</u>
Wohnsitz <u>D-69181 Leimen</u>	Residence <u>D-69181 Leimen</u>
Staatsangehörigkeit <u>deutsch</u>	Citizenship <u>german</u>
Postanschrift <u>Richard-Wagner-Str. 16</u> <u>D-69181 Leimen, Germany</u>	Post Office Address <u>Richard-Wagner-Str. 16</u> <u>D-69181 Leimen, Germany</u>
Vor- und Zuname des zweiten Miterfinders (falls zutreffend)	Full name of second joint inventor, if any
Unterschrift des zweiten Erfinders _____ Datum _____	Second Inventor's signature _____ Date _____
Wohnsitz _____	Residence _____
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(Im Falle dritter und weiterer Miterfinder sind die entsprechenden Informationen und Unterschriften hinzuzufügen.)

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